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**METHODS FOR CHOOSING
BUFFER SIZE IN
TANDEM PRODUCTION OPERATIONS**

GEORGE J. SCHLENKER

AUGUST 1983

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20. ABSTRACT (Continue on reverse side if necessary and identify by block number) This memorandum is concerned with the problem of selecting a capacity value for the buffer between two asynchronous production operations. The methodology study is motivated by shortcomings in present methods. Objectives for this study are that methods should be at least as efficient as stochastic simulation and that they should provide a means of examining both the utilization of the buffer and the productivity of the manufacturing process.		

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METHODS FOR CHOOSING
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ABSTRACT

This memorandum is concerned with the problem of selecting a capacity value for the buffer between two asynchronous production operations. The methodology study is motivated by shortcomings in present methods. Objectives for this study are that methods should be at least as efficient as stochastic simulation and that they should provide a means of examining both the utilization of the buffer and the productivity of the manufacturing process.

Two methods were developed and implemented in computer programs. Both methods use the theory of Markov processes. The first method (GS.BUF) calculates steady-state probabilities for all states of a simple production system. For certain types of processes the results of this model are exact. The results offer a good approximation for many processes. This model calculates the system productivity explicitly, providing an opportunity for economic tradeoffs between buffer capacity and other parameters. The second model (BUF.CAP) focuses on the dynamics of the filling and emptying of the buffer, under the assumption of statistical independence between states of the two operations. This model admits the possibility of several machines working in parallel at each operation.

5 August 1983

MEMORANDUM FOR RECORD

SUBJECT: Methods for Choosing Buffer Size in Tandem Production Operations

1. Reference:

a. DD 1498, HQ, US ARRCOM, DRSAR-SA, March 1983, title: Manufacturing Productivity Study.

b. Tech Report No. 82014, Menke, W. and Tran, D., ARRADCOM, November 1982, title: Simulation of Ammunition Production Lines.

2. Outline of the MFR

The following is an outline of this memorandum:

- a. Background
- b. Motivation
- c. Definitions
- d. Objectives
- e. Methodology Overview
- f. General Conclusions
- g. Derivation of Equations for GS.BUF
- h. Analytical Results of GS.BUF
- i. Conclusions Regarding GS.BUF
- j. Derivation of Equations for BUF.CAP
- k. Results of BUF.CAP
- l. Not Used
- m. Conclusions Regarding BUF.CAP
- n. Annexes -- Computer Source Programs GS.BUF and BUF.CAP

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3. Background

Over the past year I have been involved in a methodology study concerned with estimation of the capacity of production lines. This study [Ref a] has produced a general computer simulation (TANDEMT) capable of performing stochastic simulation on manufacturing systems having quite general structures. By contrast, the problem addressed by this memorandum is quite restricted in scope. One of the objectives of the manufacturing productivity study is to develop the means for understanding the significance of equipment and/or procedural changes on the productivity of a specific manufacturing line. The methods discussed here concern only the effect of buffer capacity in a simple, asynchronous system with two operations in tandem and an intervening buffer.

4. Motivation

Ref b mentions shortcomings of procedures for sizing the buffer between manufacturing operations. An ad hoc method is proposed there for generating a reasonably sized buffer, suitable for use as an input datum to a stochastic simulation of a developmental production process. The method is not claimed to be logically rigorous. The proposal yields a single number, but lacks a measure of sensitivity of the process output (or buffer performance) to buffer size. A sound approach to this problem is needed which permits the calculation of the advantage of increasing buffer capacity. A proper method should permit efficient tradeoffs to be made between buffer capacity and other process parameters. The methods presented in this memorandum possess these attributes.

5. Definitions

The simplest tandem system consists of two operations running asynchronously with an intervening buffer. In this MFR the term "simple production system" refers specifically to this kind of system. More complex systems may be viewed as arrangements of these simple systems. Additionally, the term productivity is used here somewhat restrictively. Productivity is a measure of the efficiency of a production system to produce, within imposed machine limits. As used here, productivity is defined as the ratio of the average quantity produced, in steady state, to the production of a perfect system having the same machine rates. The machines comprising the production system are viewed in a general way. They can consist of automatic hardware or of humans with simple tools or anything in between. A property of a machine is that it fails or requires adjustment by a repairman at random operating intervals. The mean time between "failures" of a machine type is a property of this machine, abbreviated MTBF. Similarly, the mean time to repair is denoted MTTR. To simplify the analysis without undue loss of generality, it is assumed that sufficient repairmen are available so that essentially no time*

* Alternatively, at least m repairmen are present, where,
 $P\{m \text{ or less repairmen busy} \} \geq 0.9.$

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is spent by machines waiting in a repair queue. This is what is meant by a well maintained system.

6. Objectives

The principal objective of this study is to develop a method for sizing the buffer in a simple production system. The method should be at least as numerically efficient as stochastic simulation. Further, the method should provide a means of examining both the utilization of the buffer and the productivity of the manufacturing process.

7. Methodology Overview

Before getting into details of the analysis, we consider the general approach to the methods of this memorandum. Both of the methods are based on the theory of Markov processes. The computer programs which implement these methods are found in the Annexes. In the first method, with program name GS.BUF, the three components of a simple production system are viewed as operating together to generate various system states. The system states are defined in terms of the admissible states of the 1st machine operation, of the buffer, and of the 2nd machine operation. As shown in Figure 1, the 1st operation is considered a single machine whose states are (a) under repair, (b) waiting for a space in the buffer to place a completed part, and (c) operating on a part. These substates are numerically coded as 0, 1, and 2 respectively. The state of the buffer is just the number of parts occupying it. The admissible buffer substates are integers from 0 to the buffer capacity. The 2nd operation is also considered a single machine whose states are coded 0, 1, and 2 for (a) under repair, (b) waiting for a part to remove from the buffer, and (c) operating on a part.

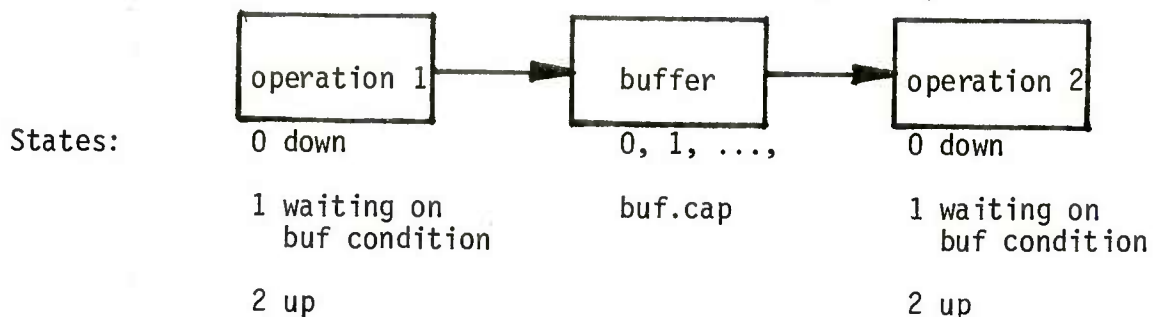


Figure 1. States of a Simple Production System

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8. A state of the system is denoted by three integers each of which characterizes one of the consecutive components of this system. Thus, state 0, 0, 1 indicates that operation 1 is under repair, that the buffer is empty, and that operation 2 is waiting for a part to operate upon. For convenience the states are labeled with a single index i . The probability that the system occupies the i th state at time t is denoted $p_i(t)$. The linear equations which functionally relate $d(p_i(t))/dt$ to the various probabilities of state occupancy, for all states, are called the Kolmogorov equations. These differential - difference equations can be simply written under the assumptions of exponentially distributed -- time to fail, time to repair, and machine operation (or service) times. (As shown later these assumptions are not too restrictive.)

9. In stochastic steady state all the derivatives are set to zero. The resultant set of linear algebraic equations is solved for the state probabilities in steady state. Certain states are identified which collectively represent interesting conditions. For example, those states in which operation 1 is waiting or operation 2 is waiting for the buffer, etc. A condition of particular interest is: operation 2 is either waiting or down for repair. The probability that this condition obtains is the fraction of time that the simple system is nonproductive. The 1's complement of this probability is, thus, system productivity. The probabilities of these state conditions can be used in tradeoff analyses to size the buffer. Economic factors can be invoked to determine the value of increased productivity versus the cost of additional buffer capacity. This is the basis of the first method for sizing the buffer.

10. A Second Approach

Another approach to buffer sizing is considered. Altho lacking in the comprehensiveness of the first approach, it does provide an approximation of the probability that a specific number of buffer spaces would be required if the first and second machine operations were not constrained to wait for a buffer condition. Unlike the first approach, this method explicitly accounts for the possible existence of several identical machines working in parallel at each machine operation. The basis of the 2nd approach is to consider the states of the first operation to be statistically independent of the states of the second. Independence is a reasonable assumption if the buffer capacity is large and if there are no common causes of machine failure. Then, the Kolmogorov equations for the first operation and for the second are independent and have the same simple form. Generally, an operation has N machines operating in parallel and possesses $N+1$ states. In this case the state value of an operation is just the number of machines operating. In a two-operation system with N_1 machines in operation 1 and N_2 machines in operation 2, there are $(N_1+1)(N_2+1)$

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system states. The method, implemented in BUF.CAP, starts with the system conditionally in each of its states in turn, and directly solves the differential - difference (Kolmogorov) equations to obtain the time-dependent average (expected) production from each of the operations, conditioned by the initial state. Because of the independence between operations, the actual numerical procedure considers each operation by itself and obtains the $N+1$ conditional trajectories -- time-dependent state probabilities and associated expected production -- storing the results. The difference in expected production from operations 1 and 2 at time t represents the expected value of the parts which would be added to the buffer, if product 1 > product 2, or would be removed from the buffer, if product 1 < product 2. Recall that this expected difference is conditional upon the initial state condition (IC). But because the IC's are random the expected production differences are actually random variables.

11. These differences are calculated for all system states at a time large enough to allow the state probabilities to approach their steady-state values (about $4 \cdot \text{MTTR}$). Then the expected differences are rank ordered from smallest to largest. The probability that the system initial state would be occupied in steady state is calculated (via a product of binomial probabilities). These probabilities are associated with each of the ordered production - difference values. These probability densities are accumulated to yield the cumulative distribution function (c.d.f.) for the production difference expected under these conditions. The program BUF.CAP displays this c.d.f. In a balanced production system the expected value of the production difference from this distribution is zero. The mean and variance of the random variable from this distribution are calculated in BUF.CAP. The variance depends upon the number of machines at each operation, the machine rates, and the MTBF's and MTTR's. To reduce the risk of causing machines to wait for the buffer, the required buffer capacity is set to the range in expected production difference plus 4 standard deviations. This somewhat arbitrary assignment provides a reasonably small risk that the calculated buffer capacity would be inadequate. A measure of risk is provided explicitly in this approach, but no measure of productivity is given here. However, using stochastic simulation I have noted that the marginal change in productivity at the calculated buffer capacity is about 0.04 percent per percent change in capacity.

12. General Conclusions

The most general conclusions concerning the above methods are presented here. Details are presented in later parts of this memorandum. In production systems having operation times which are exponentially distributed and in which repair times are exponentially distributed, a steady-state Markov model of a simple production system provides a satisfactory approach to

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sizing the buffer. A model of this type is implemented in the computer program GS.BUF. Even when the machine operating rates are constants (or nearly so), GS.BUF provides a good approach, providing that several machines are working in parallel at each operation. When constant-rate single-machine operations obtain, stochastic simulation is the preferred approach to buffer sizing. The use of GS.BUF permits one to invoke economic factors in sizing the buffer. Machine operating rate (or operating time interval), repair time, and buffer capacity are all important variables which affect productivity. If one wishes to examine the (filling and emptying) dynamics of the buffer in detail, the method implemented in BUF.CAP is suggested. This method is also useful in obtaining a point estimate of buffer capacity when a more elaborate economic tradeoff is not appropriate.

13. Derivation of Equations for GS.BUF

The general outline of the theory for GS.BUF, provided in paragraph 7 ff., will be detailed here. To make the exposition simple, consider a simple production system with a buffer having only 3 spaces. The theory of Markov processes can be applied directly to the states defined in Figure 1 providing the operating times, time intervals between failure, and repair times are all exponential random variables. Other substates could be added to accommodate other distributions of these random variables. For the present, consider only the states given in Table 1.

14. Transitions occur between these states. The state numbers in Table 1 appearing at the left in sequence are used as indices to designate the probability of state occupancy. Thus, $p_1(t)$ refers to the probability that the system is in state 1 at time t . The states to which a particular state may transition are listed in the column labeled "transition-to states." Similarly, the states from which transfers may occur are listed in the column labeled "transition-from states." Ordinarily, Markov processes may be represented diagrammatically by a graphical network with nodes as states and arcs as transitions -- the state-transition diagram. The rate parameters for the transitions are affixed to the corresponding arcs. Because of the visual complexity of the graph for this process, only a partial state-transition diagram is shown in Figure 2. In this case the first seven states are isolated, and all states connecting each of these states are shown separately. The rate parameters shown in Figure 2 -- r , λ , μ -- are indexed with a 1 or 2 to indicate the operation to which it belongs. For example, for state 1, transitions to state 2 occur at the "birth" rate λ_1 , which characterizes operation 1.

TABLE 1

DEFINITION OF STATES OF A SIMPLE PRODUCTION SYSTEM

Example with buffer capacity of 3 spaces.

State Number	State Definition*			Transition-to States	Transition-from States
	Opn 1	Buf	Opn 2		
1	0	0	1	2	2, 4
2	2	0	1	1, 6	1, 6
3	0	0	0	4, 5	4, 5
4	0	0	2	1, 3, 6	3, 6, 8
5	2	0	0	3, 6, 9	3, 6
6	2	0	2	2, 4, 5, 10	4, 5, 10
7	0	1	0	8, 9	8, 9
8	0	1	2	4, 7, 10	7, 10, 12
9	2	1	0	7, 10, 13	5, 7, 10
10	2	1	2	6, 8, 9, 14	6, 8, 9, 14
11	0	2	0	12, 13	12, 13
12	0	2	2	8, 11, 14	11, 14, 16
13	2	0	0	11, 14, 17	9, 11, 14
14	2	2	2	10, 12, 13, 18	10, 12, 13, 18
15	0	3	0	16, 17	16, 17
16	0	3	2	12, 15, 18	15, 18
17	2	3	0	15, 18, 19	13, 15, 18
18	2	3	2	14, 16, 17, 20	14, 16, 17, 20
19	1	3	0	20	17, 20
20	1	3	2	18, 19	18, 19

* For operations 1 and 2, the integers in the state definition have the following meanings:

- 0 means "down" or under repair, with a part being held.
- 1 means waiting for a buffer condition--to place a part, for operation 1, and to remove a part, for operation 2.
- 2 means "up" or working on a part.

The integer characterizing the buffer state is the number of parts in the buffer.

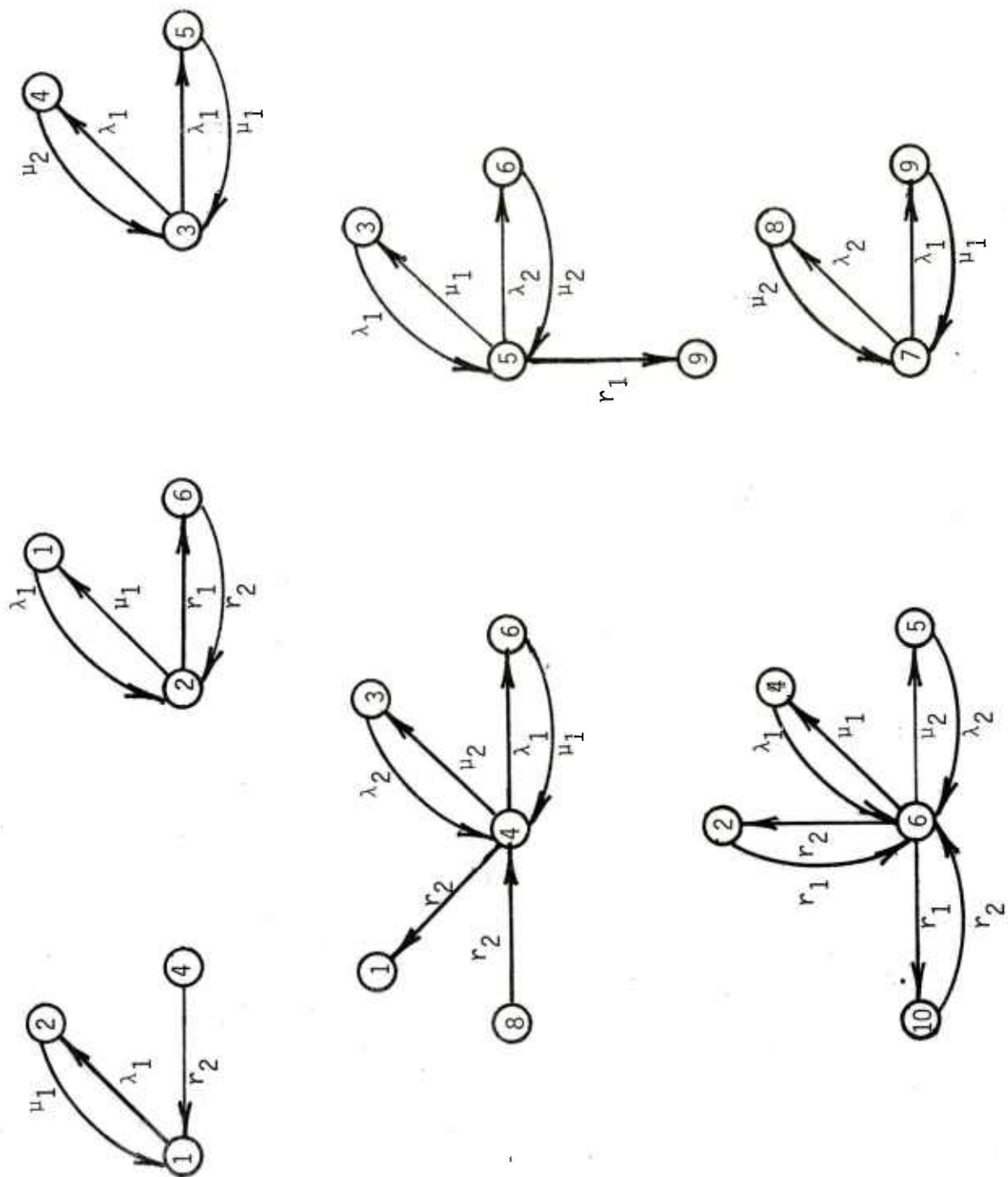


Figure 2. Partial State Transition Diagram for a Simple Production Process

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Transitions to state 1 occur from states 2 and 4. The transition from state 2 occurs at the "death" rate μ_1 , associated with operation 1. The transition from state 4 occurs at the operating rate r_2 of operation 2. The birth rate for an operation is, of course, the rate at which that operation is restored to operational condition given that it is "down". Thus,

$$\lambda_i = 1/\text{MTTR}_i, \quad i = 1, 2, \quad (1)$$

for operations 1 and 2. Similarly, the death rate for an operation is the rate at which a failure occurs, given that the operation is "up". Thus,

$$\mu_i = 1/\text{MTBF}_i, \quad i = 1, 2. \quad (2)$$

The machine rates (reciprocals of mean service time) are denoted by r_1 and r_2 , for operations 1 and 2 respectively.

15. With the aid of the state transition diagram, writing the Kolmogorov equations is a quite mechanical task. For example, as is customary in deriving transition equations, consider a small time increment h . Then, the probability that state 1 is occupied at time $t+h$ is given as

$P\{\text{state 1 at } t+h\} = P\{\text{no transition in } h \text{ occurs from state 1, given occupation at } t\}$

$*P\{\text{state 1 at } t\} + P\{\text{transition from state 2, given occupation of state 2 at } t\}$

$*P\{\text{state 2 at } t\} + P\{\text{transition from state 4, given occupation of state 4 at } t\}$

$*P\{\text{state 4 at } t\}.$

Using the abbreviated notation, this expression becomes

$$p_1(t+h) = (1-\lambda_1 h)p_1(t) + \mu_1 h p_2(t) + r_2 h p_4(t). \quad (3)$$

Then,

$$[p_1(t+h)-p_1(t)]/h = -\lambda_1 p_1(t) + \mu_1 p_2(t) + r_2 p_4(t). \quad (4)$$

Taking the limit as h approaches zero and omitting the functional dependence upon t ,

$$\dot{p}_1 = -\lambda_1 p_1 + \mu_1 p_2 + r_2 p_4. \quad (5)$$

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So much for conventional derivations! This expression can be obtained directly from the state transition diagram by writing as terms the probabilities of all states which transition to a particular state--the 1st here--on the right with their transition rates as positive coefficients. The rates of all transitions from the particular state are collected and the negative of this sum is the coefficient of the particular state. The Kolmogorov state equations for the simple system with 3-space buffer are written compactly in Table 2. The first column in this table is the index (i) of the left hand side \dot{p}_i .

16. For a simple system with a buffer capacity of 3-spaces, there are 20 states. The display of the corresponding 20 equations is awkward, but still manageable. With increasing buffer size, writing the state equations explicitly is infeasible. Fortunately, this display is not necessary. For notational simplification, let \underline{p} be a column vector of the state probabilities with ns (number of states) elements. Let B be a square matrix of coefficients with ns rows. Then, the Kolmogorov equations for a simple production system can be written, generally, as

$$\dot{\underline{p}} = B\underline{p} \quad (6)$$

Because the sum of the elements of \underline{p} is unity, B is not of full rank. Thus, equation (6) is not solved directly. In obtaining a solution, however, it is useful to generate the elements (b_{ij}) of B. Because B is a stochastic matrix, the sum of each of its column vectors is zero (reflecting the fact that each arc in the state transition diagram both leaves and enters a node). This fact is exploited to check the validity of the equations actually solved in GS.BUF.

17. In Table 1, note that the first 2 states are waiting (i.e., 1) states for operation 2, whereas the last 2 states are waiting states for operation 1. In each of the states between the 2nd and 2nd to last, a regular pattern is observed. For a given buffer state two 0 states are assigned operation 1, with operation 2 taking the values 0 and 2. Next, two 2 states are assigned for operation 1, with operation 2 again taking the values 0 and 2. This pattern of four system states is followed for each value of the buffer state. Thus, with a buffer capacity m, the number of states is

$$ns = 4(m+1)+4$$

or

$$ns = 4(m+2) \quad (7)$$

Because of the regularities in the state definition noted above, there are regularities in the equations, which permit the elements of the B matrix to be written recursively. Starting with state 11 (11th row of B),

$$b_{ij} = b_{i-4, j-4}, \quad \begin{aligned} 11 \leq i \leq ns-5, \\ i-4 \leq j \leq ns \end{aligned} \quad (8)$$

This fact greatly simplifies the process of generating the elements of the B matrix.

TABLE 2

STATE EQUATIONS FOR A SIMPLE SYSTEM
WITH 3-SPACE BUFFER

(See Table 1 for definition of states).

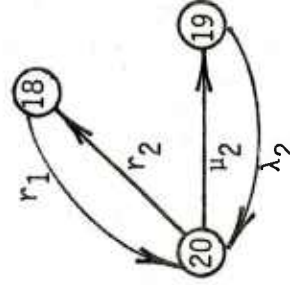
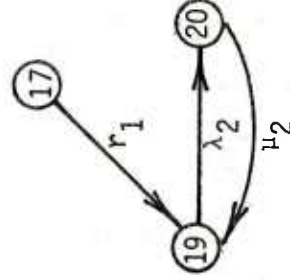
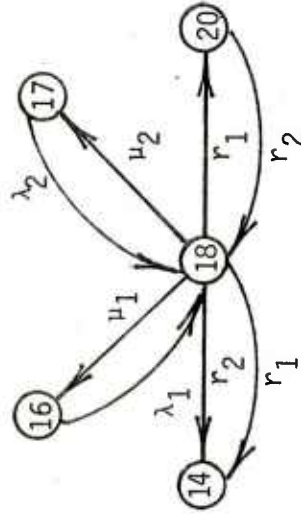
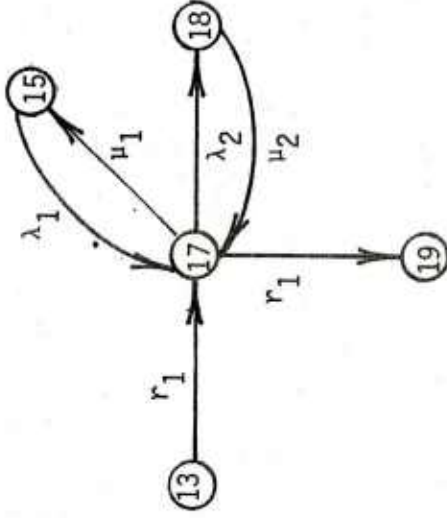
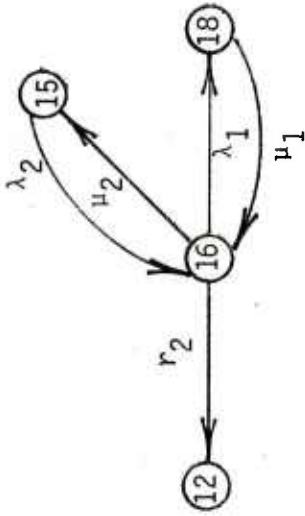
Derivative of P{State No.}	Terms in Right-Hand-Side Sum
1	$-\lambda_1 p_1, \mu_1 p_2, r_2 p_4$
2	$\lambda_1 p_1, -(\mu_1 + r_1) p_2, r_2 p_6$
3	$-(\lambda_1 + \lambda_2) p_3, \mu_2 p_4, \mu_1 p_5$
4	$\lambda_2 p_3, -(\lambda_1 + \mu_2 + r_2) p_4, \mu_1 p_6, r_2 p_8$
5	$\lambda_1 p_3, -(\lambda_2 + \mu_1 + r_1) p_5, \mu_2 p_6$
6	$r_1 p_2, \lambda_1 p_4, \lambda_2 p_5, -(\mu_1 + \mu_2 + r_1 + r_2) p_6, r_2 p_{10}$
7	$-(\lambda_1 + \lambda_2) p_7, \mu_2 p_8, \mu_1 p_9$
8	$\lambda_2 p_7, -(\lambda_1 + \mu_2 + r_2) p_8, \mu_1 p_{10}, r_2 p_{12}$
9	$r_1 p_5, \lambda_1 p_7, -(\lambda_2 + \mu_1 + r_1) p_9, \mu_2 p_{10}$
10	$r_1 p_6, \lambda_1 p_8, \lambda_2 p_9, -(\mu_1 + \mu_2 + r_1 + r_2) p_{10}, r_2 p_{14}$
11	$-(\lambda_1 + \lambda_2) p_{11}, \mu_2 p_{12}, \mu_1 p_{13}$
12	$\lambda_2 p_{11}, -(\lambda_1 + \mu_2 + r_2) p_{12}, \mu_1 p_{14}, r_2 p_{16}$
13	$r_1 p_9, \lambda_1 p_{11}, -(\lambda_2 + \mu_1 + r_1) p_{13}, \mu_2 p_{14}$
14	$r_1 p_{10}, \lambda_1 p_{12}, \lambda_2 p_{13}, -(\mu_1 + \mu_2 + r_1 + r_2) p_{14}, r_2 p_{18}$
15	$-(\lambda_1 + \lambda_2) p_{15}, \mu_2 p_{16}, \mu_1 p_{17}$
16	$\lambda_2 p_{15}, -(\lambda_1 + \mu_2 + r_2) p_{16}, \mu_1 p_{18}$

TABLE 2 (Cont)

STATE EQUATIONS FOR A SIMPLE SYSTEM
WITH 3-SPACE BUFFER

(See Table 1 for definition of states).

Derivative of $P\{\text{State No.}\}$	Terms in Right-Hand-Side Sum
17	$r_1 p_{13} , \lambda_1 p_{15} , -(\lambda_2 + \mu_1 + r_1) p_{17} , \mu_2 p_{18}$
18	$r_1 p_{14} , \lambda_1 p_{16} , \lambda_2 p_{17} , -(\mu_1 + \mu_2 + r_1 + r_2) p_{18} , r_2 p_{20}$
19	$r_1 p_{17} , -\lambda_2 p_{19} , \mu_2 p_{20}$
20	$r_1 p_{18} , \lambda_2 p_{19} , -(\mu_2 + r_2) p_{20}$



NOTE: For a general system with ns states, the above state numbers correspond as follow:
 $ns = 20$, $ns - 1 = 19$, etc.

Figure 3. State Transition Diagram of the Last Five States for a Simple Production Process

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18. The last 5 system states, which contain waiting states for operation 1, must be written explicitly. For the general case, these are:

$$\dot{p}_{ns-4} = \lambda_2 p_{ns-5} - (\lambda_1 + \mu_2 + r_2) p_{ns-4} + \mu_1 p_{ns-2} \quad (9)$$

$$\dot{p}_{ns-3} = r_1 p_{ns-7} + \lambda_1 p_{ns-5} - (\lambda_2 + \mu_1 + r_1) p_{ns-3} + \mu_2 p_{ns-2} \quad (10)$$

$$\dot{p}_{ns-2} = r_1 p_{ns-6} + \lambda_1 p_{ns-4} + \lambda_2 p_{ns-3} - (\mu_1 + \mu_2 + r_1 + r_2) p_{ns-2} + r_2 p_{ns} \quad (11)$$

$$\dot{p}_{ns-1} = r_1 p_{ns-3} - \lambda_2 p_{ns-1} + \mu_2 p_{ns} \quad (12)$$

$$\dot{p}_{ns} = r_1 p_{ns-2} + \lambda_2 p_{ns-1} - (\mu_2 + r_2) p_{ns} \quad (13)$$

The state transition diagram for the last five states is found in Figure 3.

19. To evaluate the system under stochastic steady state, the \dot{p} vector is set to zero.

Then,

$$B \underline{p} = \underline{0} \quad (14)$$

As noted above, the resultant set of equations contains one superfluous equation, since the state probabilities must sum to 1. To remedy this situation, (14) is converted to the linear matrix-vector equation

$$A \underline{x} = \underline{c}$$

with

$$A \text{ (nxn)}, \quad \underline{x} \text{ (nx1)}, \text{ and } \underline{c} \text{ (nx1)},$$

where

$$n = ns-1$$

This is done by assigning

$$p_1 = 1 - \sum_{k=2}^{ns} p_k \quad (16)$$

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The elements of the A matrix and \underline{c} vector in (15) are obtained via the transformations

$$b'_{ij} = b_{ij} - b_{i1} \quad , \quad 2 \leq i \leq ns$$

and,

$$c_{i-1} = b_{i1} \quad , \quad 2 \leq i \leq ns$$

$$a_{ij} = b'_{i+1, j+1} \quad , \quad 1 \leq i, j \leq n \quad . \quad (17)$$

Note that the first scalar equation in equation (14) is deleted in forming (15). The A matrix is of full rank, so \underline{x} may be obtained by

$$\underline{x} = A^{-1} \underline{c} \quad . \quad (18)$$

Then,

$$p_{i+1} = x_i \quad , \quad 1 \leq i \leq n \quad . \quad (19)$$

Finally, p_1 is obtained from (16).

20. The probabilities of buffer occupancy are obtained from the state probability vector \underline{p} (or, alternatively from \underline{x}) by

$$P\{\text{buffer is empty}\} = \sum_{k=1}^6 p_k \quad . \quad (20a)$$

$$P\{j \text{ parts in buffer}\} = \sum_{k=1}^4 p_{k+4j+2} \quad , \quad 1 \leq j \leq m-1 \quad . \quad (20b)$$

$$P\{\text{buffer is full}\} = (\text{by definition})$$

$$P\{m \text{ parts in buffer}\} = \sum_{k=1}^6 p_{k+4m+2} \quad . \quad (20c)$$

The probability that operation 2 must wait for a buffer condition is just the sum of the first two state probabilities:

$$P\{\text{operation 2 must wait}\} = p_1 + p_2 \quad . \quad (21)$$

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The probability that operation 1 must wait for a buffer condition is the sum of the last two state probabilities:

$$P\{\text{operation 1 must wait}\} = p_n + p_{n+1} \quad (22)$$

The system productivity depends upon the probability that operation 2 is down. This last probability is given by

$$P\{\text{operation 2 is down}\} = \sum_{k=1}^{(ns-2)/2} p_{2k+1} \quad (23)$$

Finally, the system productivity is

$$1 - P\{\text{operation 2 is down}\} - P\{\text{operation 2 must wait}\} \quad (24)$$

If one wants the steady-state production rate of the simple system, this value is obtained from the productivity by multiplying times r_2 .

21. Analytic Results of GS.BUF

Calculations were made using GS.BUF to compare results with stochastic simulation and to do a sensitivity analysis of certain parameters. This computational experience provided timing estimates on our PRIME 550 mini-computer. All calculations were made using double-precision arithmetic. The stochastic simulation used for comparison was a very simple implementation of TANDEMT. In all cases studied the time between failures and the time to repair were exponential random variables. Simulation runs were 40 24-hour days, starting with an empty system. Runs were made for instances in which the machine operating (or service) time is exponential and in which it is constant. The runs with constant service time are used to test the applicability of the Markov model in GS.BUF to a quite different model. Parenthetically, I note that other Markov models can be created to approximate the constant-rate case. By defining many substates to describe a machine operation, it is possible to describe a random service time whose coefficient of variation is quite small (if not zero). In fact, I constructed a specific case of such a model using 3 substates. The productivity calculated with that model is in better agreement with simulated results than is the case for GS.BUF. However, due to the large number of system states produced by this method, computational efficiency is poor, for a typical buffer size. Consequently, that approach was abandoned. It is recommended for single-machine operations having constant rates, that stochastic simulation be used to estimate system productivity. This approach also has the advantage of modeling other-than well maintained systems, which were considered here.

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22. Results of GS.BUF in which the operations have a common set of parameters are shown on the first page of Table 3. Buffer capacity is treated as a parameter in the comparison between calculated and simulated results. This comparison indicates complete agreement within expected statistical variation. For simulations of this length the estimated standard error of the productivity estimate is about 0.11 to 0.14 for these examples. The typically somewhat larger estimated probability that the buffer is empty is a reflection that the simulation was not started in the steady state. However, due to the simulation length, this effect is small. On the second page of Table 3, a comparison is made in which the parameters of operation 1 are not the same as those of operation 2. Again, agreement is excellent. Whenever one is designing a buffer, a balanced pair of operations should be considered. When balanced, the operation thruput is a constant equal to the product of the machine rate, the machine availability, and the number of machines. Note that these operations are balanced.

23. Calculated and simulated productivities for an expanded set of buffer capacities are shown in Table 4. This sort of analysis can be used in making economic tradeoffs when choosing a buffer capacity. The results of TANDEMT are shown in Table 5. It is noted that the system productivity is always greater when the machine service (operating) times are constants than when they are exponential random variables. However, note that the difference (and relative difference) in productivities in these two instances diminishes as the buffer capacity increases. This fact suggests that regardless of the distribution of service times, GS.BUF may be a practical procedure to use for productivity estimation (and tradeoff) when the buffer capacity is large. Figure 4 shows the probability density functions of buffer state occupancy for cases in which the service times are constant and exponential random variables. The U-shaped densities are typical. Note that much larger probabilities of being at the buffer extremes is exhibited by a constant-rate system.

24. The sensitivity analysis using GS.BUF considers the effect of the following three parameters on system productivity: buffer capacity, MTTR, and machine rate. While not large, the ranges of these parameters are representative of many ammunition production operations. Several inferences can be drawn from the results shown in Table 6. Only one is mentioned at this point. Suppose the machine rate is given and the buffer size chosen to yield a particular productivity or, alternatively, chosen so that the marginal cost of additional buffer space just equals the value of additional production from a system with the incrementally larger buffer. If, later, a greater machine rate is available via, possibly, machine substitution, one can expect a productivity decrease if the buffer is not enlarged. Remember that productivity is a measure of production efficiency. Thus, doubling the machine rate -- with buffer fixed -- will increase the production rate, but will not double it.

TABLE 3

COMPARISON OF THE PROBABILITIES OF BUFFER
OCCUPANCY: CALCULATED VERSUS SIMULATED

Parameters:

Two single-machine operations in tandem.

Operation (service) times are exponential.

Repair times are exponential.

MTBF = 100 minutes

MTTR = 25 minutes

Average operating rate = 1 part/min.

Buffer Capacity	Buffer Status	State Probabilities	
		Calculated	Simulated
3	0	0.379	0.385
	1	0.121	0.125
	2	0.121	0.126
	3	0.379	0.364
		System Productivity*	
		0.595	0.599
5	0	0.320	0.312
	1	0.091	0.094
	2	0.089	0.095
	3	0.089	0.094
	4	0.091	0.094
	5	0.320	0.311
		System Productivity*	
		0.620	0.630
10	0	0.250	0.247
	1	0.059	0.061
	2	0.056	0.059
	3	0.055	0.056
	4	0.054	0.054
	5	0.053	0.056
	6	0.054	0.054
	7	0.055	0.056
	8	0.056	0.057
	9	0.059	0.060
	10	0.250	0.240
		System Productivity*	
		0.650	0.640

TABLE 3 (Cont)

COMPARISON OF THE PROBABILITIES OF BUFFER
OCCUPANCY: CALCULATED VERSUS SIMULATED

Parameters:

Two single-machine operations in tandem.

Operation (service) times are exponential.

Repair times are exponential.

MTBF = 200 min (1st opn), = 100 min (2nd opn).

MTTR = 25 min (1st opn), = 12.5 min (2nd opn).

Average operating rate = 1/min.

Buffer Capacity	Buffer States	State Probabilities	
		Calculated	Simulated
3	0	0.357	0.368
	1	0.140	0.143
	2	0.141	0.144
	3	0.362	0.345
		System Productivity*	
		0.695	0.696
5	0	0.287	0.291
	1	0.103	0.105
	2	0.104	0.106
	3	0.105	0.106
	4	0.108	0.109
	5	0.294	0.283
		System Productivity*	
		0.727	0.731

* System productivity is the ratio of the average production achieved to the maximum steady-state production from a system of perfect machines operating at the same rates.

TABLE 4

COMPARISON OF CALCULATED WITH
SIMULATED PRODUCTIVITY ESTIMATES FOR
A SIMPLE PRODUCTION SYSTEM*

Buffer Capacity	Average Productivity	
	Calculated	Simulated**
3	0.595	0.599
5	0.620	0.630
10	0.650	0.640
20	0.680	0.684
40	0.710	0.713

* Parameters:

A single machine at each of two operations.

Exponential operation times.

Exponential repair times.

Average operating rate 1 part/minute.

Common MTBF = 100 minutes.

Common MTTR = 25 minutes.

** The standard error is about 0.011, based on a 40 day simulation.

TABLE 5

COMPARISON OF PRODUCTIVITY ESTIMATES FOR
SIMPLE PRODUCTION SYSTEMS* HAVING
CONSTANT VERSUS EXPONENTIAL OPERATING TIMES

Tabulated productivity is the ratio of expected production to the maximum production from a perfect system operating at the same machine rate.

Buffer Capacity	<u>Simulated Operating Times</u>	
	Constant	Exponential
3	0.674	0.599
5	0.679	0.630
10	0.690	0.640
20	0.709	0.684
40	0.724	0.713

* Parameters:
Single machine operations.
Average operating rate 1/min.
Repair times are exponential.
MTBF = 100 minutes
MTTR = 25 minutes



TABLE 6

CALCULATED PRODUCTIVITY OF A SIMPLE PRODUCTION
SYSTEM* AS A FUNCTION OF SEVERAL VARIABLES

The tabulated productivity is the ratio of average production achieved to the maximum steady-state production from a perfect system operating at the same machine rate.

Buffer Capacity	MTTR (minutes)	Machine Rate (parts/min)	
		1	2
10	12.5	0.770	0.761
	25.0	0.650	0.642
20	12.5	0.805	0.792
	25.0	0.680	0.666
40	12.5	0.835	0.819
	25.0	0.710	0.690

* System parameters:

Two single-machine operations in tandem.
Capacity of intermediate buffer is a parameter.
Common machine rate is a parameter.
Common MTTR is a parameter.
Exponential operation (service) times.
Exponential repair times.
Common MTBF = 100 minutes.

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25. The system model of GS.BUF assumes single-machine operations. However, multi-machine operations can be approximated by this model by scaling the machine rate to Nr , where N is the number of machines in the operation, each having rate r . This approximation exploits a theorem of random processes on the pooling of many component processes. Cox* states that when many independent process events are pooled, the pooled process is approximately Poisson, i.e., the time between events is approximately exponential. This is true irrespective of the distributions of the component processes. In application to multi-machine operations, each machine's output gets pooled for that operation. Thus, the times between unit production events are approximately exponential random variables when the number of machines is large. Surprisingly, the number working in parallel at an operation does not need to be more than about 4 to yield a good approximation of the probability density function (p.d.f.) for buffer occupancy. This point is illustrated by the results in Table 7. Two multi-machine cases were simulated: 4 machines per operation working at machine rate 1/4 parts per minute and 5 machines working in parallel at machine rate 1/5 parts per minute. For comparison is shown the p.d.f. of buffer occupancy calculated with GS.BUF having a machine rate of 1 part per minute. One observes that the p.d.f.'s in these instances are nearly the same. Agreement between calculated and simulated productivities is not as good, however. The simulated productivities for the two examples are both 0.87 whereas the calculated productivity is 0.81.

26. Conclusions Regarding GS.BUF

Specific conclusions for the method of GS.BUF are summarized here.

(a) When a simple production system satisfies the assumption of exponentially distributed service times, the results of GS.BUF are exact. This point has been verified with simulation. (b) With stochastic steady state, the probability density function (p.d.f.) of buffer occupancy is U-shaped. For operations having a common set of operating characteristics, the above p.d.f. is symmetric with respect to $0.5 \cdot$ buffer capacity spaces. A prominent positive jump in the p.d.f. is observed at the first and last states. For a given buffer capacity the jump is more pronounced when the operating times are constant than when they are both exponential. This implies that relatively more time is spent at extreme states when the operating times are constant than when they are random. Nevertheless, the productivity of a system with constant operating rate is greater than that of a system with a random rate of the same mean value, other things being the same.

* Page 77 ff. Cox, D.R. Renewal Theory, London, distributed by Barnes and Noble, c. 1962.

TABLE 7

COMPARISON OF AN ANALYTIC APPROXIMATION* WITH SIMULATED
PROBABILITY DISTRIBUTIONS OF BUFFER OCCUPANCY FOR
SYSTEMS HAVING MULTIPLE MACHINES PER OPERATION

* Parameters of the Analytic Model:
Two single-machine operations in tandem.
Operation times are exponential.
Repair times are exponential.
MTBF for each machine = 100 minutes.
MTTR for each machine = 12.5 minutes.
Average operating rate = 1 part/min.
Buffer capacity = 20 spaces.

Buffer State	Calc. pdf	Simulated ⁺ pdf with	
		4 mach at rate 1/4	5 mach at rate 1/5
0	0.141	0.131	0.135
1	0.044	0.037	0.042
2	0.042	0.035	0.043
3	0.040	0.036	0.044
4	0.038	0.037	0.043
5	0.037	0.037	0.041
6	0.036	0.041	0.036
7	0.036	0.042	0.036
8	0.035	0.041	0.038
9	0.035	0.041	0.036
10	0.035	0.039	0.033
11	0.035	0.035	0.034
12	0.035	0.039	0.034
13	0.036	0.040	0.035
14	0.036	0.040	0.037
15	0.037	0.040	0.039
16	0.038	0.037	0.037
17	0.040	0.035	0.036
18	0.042	0.037	0.040
19	0.044	0.037	0.043
20	0.141	0.143	0.138

+ Simulated machine operating rates are constants. The MTBF and MTTR parameters of the simulation are as given above.

(c) A sensitivity analysis of system parameters was performed using GS.BUF. All the following are shown to be important: Machine operating rate (or service time), repair time, and buffer capacity. (d) If the buffer capacity remains constant and the machine rate is doubled, a productivity decrease is experienced -- even tho the production rate of the new configuration is greater. This loss of productivity is more pronounced for a system with a large buffer than with a system whose buffer is inadequate. (e) Doubling the production rate and concurrently doubling buffer capacity increases the system productivity. Thus, the buffer size need not be doubled to preserve productivity if the operating rate is doubled. (f) At any reasonable buffer size, halving the MTTR produces a greater improvement in productivity than that achieved by doubling buffer capacity. (g) For systems having constant machine operating rates, GS.BUF can be used to approximate the system, only if many machines are working in parallel at each operation. The probability distribution of buffer state occupancy is approximately correct in this case. If only one fixed-rate machine exists per operation, it is recommended that stochastic simulation be used to estimate the system productivity.

27. Derivation of Equations for BUF.CAP

To motivate subsequent discussion, consider the state dynamics of a one-machine system. This system is regarded as operating independently of other buffers and production operations. In this case there are only two states -- down (0) and up (1). Markov transitions from down to up occur at the birth rate λ , and transitions from up to down occur at the death rate μ . Thus, the equations for the probabilities of state occupancy are

$$\dot{p}_0(t) = -\lambda p_0 + \mu p_1 \quad (25a)$$

$$\dot{p}_1(t) = \lambda p_0 - \mu p_1 \quad (25b)$$

These equations can be solved directly or by using Laplace transforms and inverting.

$$p_0(t) = p_0(0)e^{-(\lambda+\mu)t} + \frac{\mu}{\lambda+\mu} (1-e^{-(\lambda+\mu)t}) \quad (26a)$$

$$p_1(t) = \frac{\lambda}{\lambda+\mu} (1-e^{-(\lambda+\mu)t}) + p_1(0)e^{-(\lambda+\mu)t} \quad (26b)$$

Note that the steady-state availability,

$p_1(\infty)$, or

$$A = \frac{\lambda}{\lambda+\mu} \quad (27)$$

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Using the definitions of λ and μ , equations (1, 2),

$$A = \frac{MTBF}{MTBF+MTTR} \quad (28)$$

The (mathematically) expected production from this operation over the time interval $(0, t)$ is given by

$$c(t) = r \int_0^t p_1(x) dx, \quad (29)$$

where r is the machine operating rate.

From (26b) and (29),

$$\begin{aligned} \frac{c(t)}{r} &= \frac{\lambda}{\lambda+\mu} t + \frac{\lambda}{(\lambda+\mu)^2} e^{-(\lambda+\mu)t} \\ &- \frac{\lambda}{(\lambda+\mu)^2} + \frac{p_1(0)}{\lambda+\mu} (1-e^{-(\lambda+\mu)t}) \quad (30) \end{aligned}$$

28. Consider two single-machine operations in series with a common set of operational parameters. When first observed, let the first operation be in state 1 and the second operation be in state 2. Call this initial condition IC1. Stated mathematically, let

$$p_1(0) = 1 \quad \text{for operation 1}$$

and

$$p_1(0) = 0 \quad \text{for operation 2.}$$

Then, the difference in expected production of these operations as t becomes large can be obtained from (30) as $E[\text{production difference, given IC1}] = r/(\lambda+\mu)$. (31)
(Notationally, E is the expected value operator.) If IC1 obtains, one expects to add the above production quantity to the contents of a buffer between the operations, given the assumed freedom from buffer constraints. Clearly, if the initial operation states had been the 1's complements of the above (IC2), a quantity of production would have been removed from the buffer equal to (31). This again assumes that at least that much was present initially. The IC's considered are the extreme conditions for this example. Therefore, at reasonably low risk of buffer inadequacy one might suggest using the range of $E[\text{production difference}]$ over extreme IC's for sizing the buffer. In this case

$$\begin{aligned} &E[\text{production difference, IC1}] - E[\text{production difference, IC2}] \\ &= 2r/(\lambda+\mu) \quad (32) \end{aligned}$$

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However, this approach would ignore some obvious facts. First, we are dealing with conditional expected values. Thus, variation from these averages must occur. Secondly, we ignore the fact that the occupancy of the buffer at the commencement of delivery of additional product under IC1 may exceed $r/(\lambda+\mu)$. To reduce the risk associated with these contingencies, one can identify the variation in expected buffer state change. Specifically, one can calculate the standard deviation (SD) of this random variable. Then, some number of SD's can be added to the range in (32) to produce a buffer capacity requirement. I have somewhat arbitrarily chosen 4 SD's as a reasonably cautious number. The probability distribution of expected production difference is calculated by enumerating all possible IC's. In this example there are four -- (0, 0), (0, 1), (1, 0), and (1, 1). The probability that the system of two operations exists in one of these states is simply the product of the probability that operation 1 is in its given state and the probability that operation 2 is in its state. The probability that a one-machine operation is up (1) is A, and down (0) is (1-A). In general, in a N-machine operation

$$\begin{aligned} &P\{k \text{ machines are operating}\} \\ &= \binom{N}{k} A^k (1-A)^{N-k} . \end{aligned} \quad (33)$$

29. Clearly, the expected production difference between operations started in the same state is zero. Hence, in this example the (0, 0) and (1, 1) IC's yield $E[\text{production difference}] = 0$. Thus, the probability density of the $E[\text{production difference}]$ for this example is the following: $-r/(\lambda+\mu)$ with probability $A(1-A)$, 0 with probability $A^2+(1-A)^2$, and $r/(\lambda+\mu)$ with probability $A(1-A)$. Since the mean value of this variable is zero, the variance is given as

$$[2r^2/(\lambda+\mu)^2]A(1-A) \quad (34)$$

or, from (27), the variance of $E[\text{production difference}]$ is

$$\frac{2r^2\lambda\mu}{(\lambda+\mu)^4} , \quad (35)$$

and the SD is

$$\frac{(2\lambda\mu)^{1/2}r}{(\lambda+\mu)^2} . \quad (36)$$

As a specific numerical example, let $r=1$ part per minute, MTBF = 100 minutes. MTTR = 25 minutes. Then, from (32), the expected range in production differences is 40 parts and, from (36), the standard deviation 11.3 parts.

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This yields a required capacity of 85 spaces. To examine the marginal advantage in productivity for this buffer size, I ran GS.BUF at a buffer capacity of 85 and, again at 80. The productivities in these instances are 0.7426 and 0.7402, respectively. This implies a productivity change of about 0.04% per percent change in buffer capacity.

30. A Two-Machine Operation

As in the above example, it is not difficult to obtain an analytic solution to the time-dependent Markov model of a 2-machine production system. With two machines the system states are 0, 1, and 2. However, since

$$p_0(t) = 1 - p_1(t) - p_2(t) , \quad (37)$$

only two equations are required to describe the system. Dropping the notation for explicit time dependence,

$$\dot{p}_1 = 2\lambda(1-p_1-p_2) - (\lambda+\mu)p_1 + 2\mu p_2 \quad (38a)$$

$$\dot{p}_2 = \lambda p_1 - 2\mu p_2 . \quad (38b)$$

After some manipulation of these equations and using Laplace transforms, one obtains

$$p_1(t) = \frac{2\lambda\mu}{(\lambda+\mu)^2} + Ae^{-(\lambda+\mu)t} + Be^{-2(\lambda+\mu)t} , \quad (39a)$$

with

$$A = \frac{a_0 - a_1(\lambda+\mu)}{\lambda+\mu} - \frac{4\lambda\mu}{(\lambda+\mu)^2} \quad (39b)$$

$$B = - \frac{a_0 - 2a_1(\lambda+\mu)}{\lambda+\mu} + \frac{2\lambda\mu}{(\lambda+\mu)^2} \quad (39c)$$

$$a_0 = \dot{p}_1(0) + 3(\lambda+\mu)p_1(0)$$

or

$$a_0 = 2\lambda p_0(0) + 2(\lambda+\mu)p_1(0) + 2\mu p_2(0) \quad (39d)$$

$$a_1 = p_1(0) . \quad (39e)$$

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$$p_2(t) = p_2(0)e^{-2\mu t} + \frac{\lambda^2}{(\lambda+\mu)^2} - \frac{\lambda^2}{(\lambda+\mu)^2} e^{-2\mu t} + \frac{A}{\lambda-\mu} (e^{-2\mu t} - e^{-(\lambda+\mu)t}) + \frac{B}{2} (e^{-2\mu t} - e^{-2(\lambda+\mu)t}) . \quad (40)$$

Steady-state results are obtained from (37), (39), and (40) by allowing t to approach ∞ .

$$p_0(\infty) = \frac{\mu^2}{(\lambda+\mu)^2} \quad (41a)$$

$$p_1(\infty) = \frac{2\lambda\mu}{(\lambda+\mu)^2} \quad (41b)$$

$$p_2(\infty) = \frac{\lambda^2}{(\lambda+\mu)^2} . \quad (41c)$$

The expected number of machines operating in steady-state in this case is

$$E[\text{number operating}] = p_1(\infty) + 2p_2(\infty) , \quad (42a)$$

$$= 2\lambda/(\lambda+\mu) . \quad (42b)$$

Since the number of machines, N , is 2 in this example, as anticipated,

$$E[\text{number operating}] = NA . \quad (43)$$

The expected number of machines operating equals the max number times the intrinsic availability only if the system is well maintained. In writing the system equations (38a and b), it is assumed that if N machines are down for repair, N repairmen are working. Thus, the transition rate from the 0 state to the 1 state is given as $N\lambda$. If there were only M repairmen available for machine repairs, where $M < N$, the largest transition rate from a lower to a higher state would be $M\lambda$. In the latter case, the system state probabilities would depend upon both the number of machines per operation and the number of repairmen. To avoid this complication, it is assumed that the system is well maintained. This assumption is not as restrictive as it may seem. Approximately correct results are obtained for far weaker conditions, as will be shown later.

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31. An N-Machine Operation

The state transition diagram for the general case is shown in Figure 5. Note that the transition from the j th to the $j+1$ st states occurs at the rate $N-j$ (machines down) times λ . The downward transition from the k th to the $k-1$ st state occurs at the rate k (more to fail) times the unit death rate μ . As was done for the previous Markov models, the Kolmogorov equations for this model can be written by inspection from the state transition diagram. Since this process is quite straightforward, it is not repeated here. The general result, with the deletion of the zero (or null) state is the familiar form

$$\dot{\underline{p}}(t) = A\underline{p}(t) + \underline{c} \quad , \quad (44)$$

with $\underline{c}' = [N\lambda, 0, 0, \dots, 0]$,

where \underline{p} and \underline{c} are $(N \times 1)$ and where A is $(N \times N)$.

32. The solution procedure employs numerical integration using the rectangle rule with time step h :

$$\underline{p}(t+h) = \underline{p}(t) + h\dot{\underline{p}}(t) \quad , \quad (45)$$

with $\dot{\underline{p}}(t)$ given by (44). For a small time step (h), this procedure was found to be slightly faster (and easier) to implement than to use a double step ($2h$) with Euler's rule with a predictor and a corrector. A time step h of 0.1 minute is used in BUF.CAP. Notationally, let the conditional expected value of the output from an operation at time t (from the IC) be denoted $\bar{x}(t)$. By definition,

$$\bar{x}(t) = \int_0^t r \sum_{k=1}^N k p_k(t) dt \quad , \quad (45)$$

where r is the unit machine rate for this operation. The integrand in (45) is the average rate of production -- \underline{dx}/dt -- from this machine operation. This derivative is saved at time intervals of Δ for optional printing. Call the numerical approximation $\bar{x}(i\Delta)$, with integer i . In performing the numerical integration to calculate $\bar{x}(t)$, it is unnecessary to use a step as small as h . To yield about the same relative precision as obtained in calculating $\underline{p}(t)$, one can use a step Δ -- called DELTAT in BUF.CAP -- of 0.5 minute with Euler's rule:

$$\bar{x}(t+\Delta) \approx 0.5\Delta[\dot{\bar{x}}(t+\Delta) + \dot{\bar{x}}(t)] \quad . \quad (46)$$

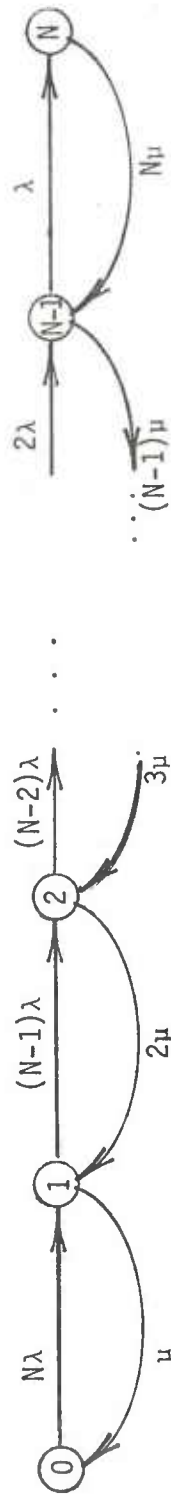


Figure 5. State Transition Diagram for a Well Maintained N-Machine Operation

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Denote by \bar{x}_{ij} the conditional expected output of operation i with initial condition (IC) j . Then, the difference in expected production values is calculated at time t_{\max} :

$$\overline{\Delta x_j} = \bar{x}_{1j}(t_{\max}) - \bar{x}_{2j}(t_{\max}) \quad . \quad (47)$$

This difference is essentially constant beyond $t_{\max} = 2\max(\text{MTTR}_1, \text{MTTR}_2)$. These expected production differences are treated in BUF.CAP in the manner described for the single-machine case (pgf. 28 ff). The probability distribution of $\overline{\Delta x_j}$ is displayed in order to facilitate the choice of buffer capacity on the basis of risk that a particular value is inadequate.

33. Sample Output from BUF.CAP

A sample run using the program BUF.CAP is shown in Table 8. The program input values are repeated at the top of the page. The output represents an abbreviated form of the available outputs. If the user chooses to display the trajectories of the conditional state probability vectors and cumulative production he may.

34. Stochastic simulation is used in the process of examining the advantage of increasing the buffer capacity beyond that required by BUF.CAP. Examples of the c.d.f.'s of buffer occupancy for several two-operation systems are shown in Figures 6 and 7. The max operation rate of both operations in each system is the same. In all cases shown here the machine service (or operation) time is constant. The probability distributions of buffer occupancy for two balanced systems are compared in Figure 6. One can observe that increasing the buffer capacity from 40 to 100 spaces, for the parameters shown, has the effect of significantly reducing the risk of encountering a full buffer. But, the probability of an empty buffer is nearly the same in both instances. Several comparisons are made in Figure 7. In these comparisons the buffer capacity is fixed at 40 spaces, and the operational rate is a constant 1 part/minute. An effect on the c.d.f. of buffer occupancy occurs when the number of machines per operation increases. The effect of this increase is to reduce the probabilities associated with the extreme states of the buffer. This phenomenon is also shown in the output of BUF.CAP. Another observation of interest can be drawn from Figure 7. Thruout this study all analyses have been conducted assuming a well maintained system. As indicated, this implies at least as many repairmen as machines. For comparison, one simulation run was made with a system of 6 machines -- 4, in operation 1 and 2, in operation 2-- and with only one repairman. The c.d.f. of buffer occupancy for this case is shown in Figure 7. A rather small difference exists between the c.d.f. for this case and the c.d.f. for a comparable, well-maintained system. The reason

TABLE 8

SAMPLE OUTPUT FROM PROGRAM BUF.CAP WITH TERMINAL DIALOG DELETED

MACHINE INPUT DATA FOR BUFFER CAPACITY CALCULATION

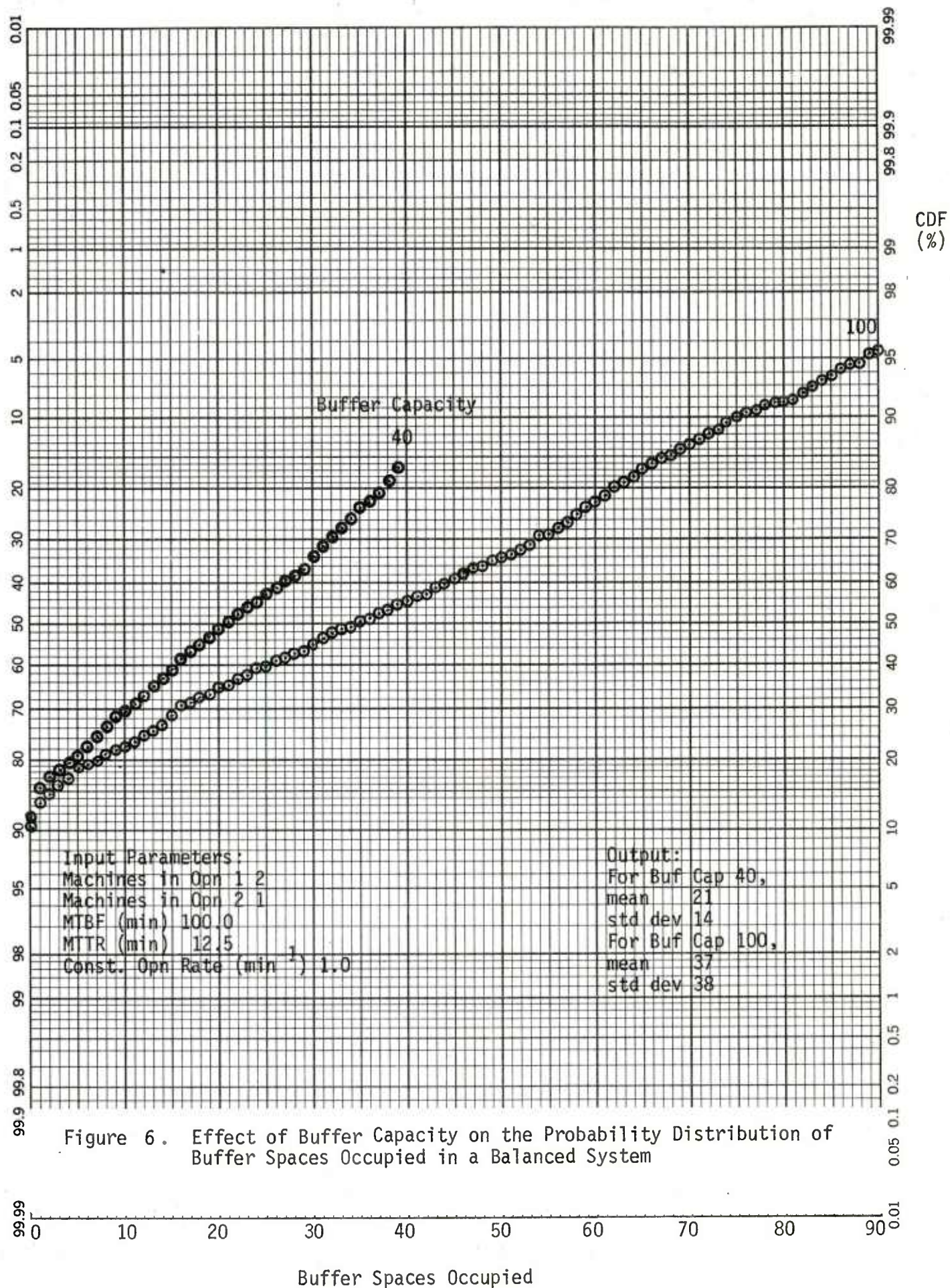
OPERATION 1		OPERATION 2	
NO MACHINES	4	4	
MACH RATE	0.25	0.25	PARTS/MIN
MTBF	100.00	100.00	MINUTES
MTTR	12.50	12.50	MINUTES
AVAILABILITY	88.89	88.89	PERCENT
REPAIR TIME (MINUTES)		50.0	ASSOCIATED STAT CONFIDENCE 0.982

EXPECTED BUFFER REQUIREMENTS ORDERED OVER ALL SYSTEM STATES
BASED ON A REPAIR TIME LAG OF 50.0 MINUTES

ORDER INDEX	BUFFER SPACES	DIFFER SPACES	PROB DENS	CUML PROB	OPN 1 STATE	OPN 2 STATE
1	-10.9	0.0	0.00001	0.00001	0	4
2	-8.2	2.7	0.0030	0.0031	1	4
3	-8.2	2.7	0.0000	0.0032	0	3
4	-5.5	5.5	0.0000	0.0032	0	2
5	-5.5	5.5	0.0365	0.0397	2	4
6	-5.5	5.5	0.0015	0.0413	1	3
7	-2.7	8.2	0.1949	0.2361	3	4
8	-2.7	8.2	0.0000	0.2361	0	1
9	-2.7	8.2	0.0003	0.2364	1	2
10	-2.7	8.2	0.0183	0.2547	2	3
11	0.0	10.9	0.0000	0.2547	1	1
12	0.0	10.9	0.0000	0.2547	0	0
13	0.0	10.9	0.0974	0.3521	3	3
14	0.0	10.9	0.0034	0.3556	2	2
15	0.0	10.9	0.3897	0.7453	4	4
16	2.7	13.7	0.0183	0.7636	3	2
17	2.7	13.7	0.0003	0.7639	2	1
18	2.7	13.7	0.0000	0.7639	1	0
19	2.7	13.7	0.1949	0.9587	4	3
20	5.5	16.4	0.0015	0.9603	3	1
21	5.5	16.4	0.0365	0.9968	4	2
22	5.5	16.4	0.0000	0.9968	2	0
23	8.2	19.1	0.0000	0.9969	3	0
24	8.2	19.1	0.0030	0.9999	4	1
25	10.9	21.9	0.0001	1.0000	4	0

AVERAGE BUFFER CHANGE ----- -0.00
 STD DEV BUFFER CHANGE ----- 2.43

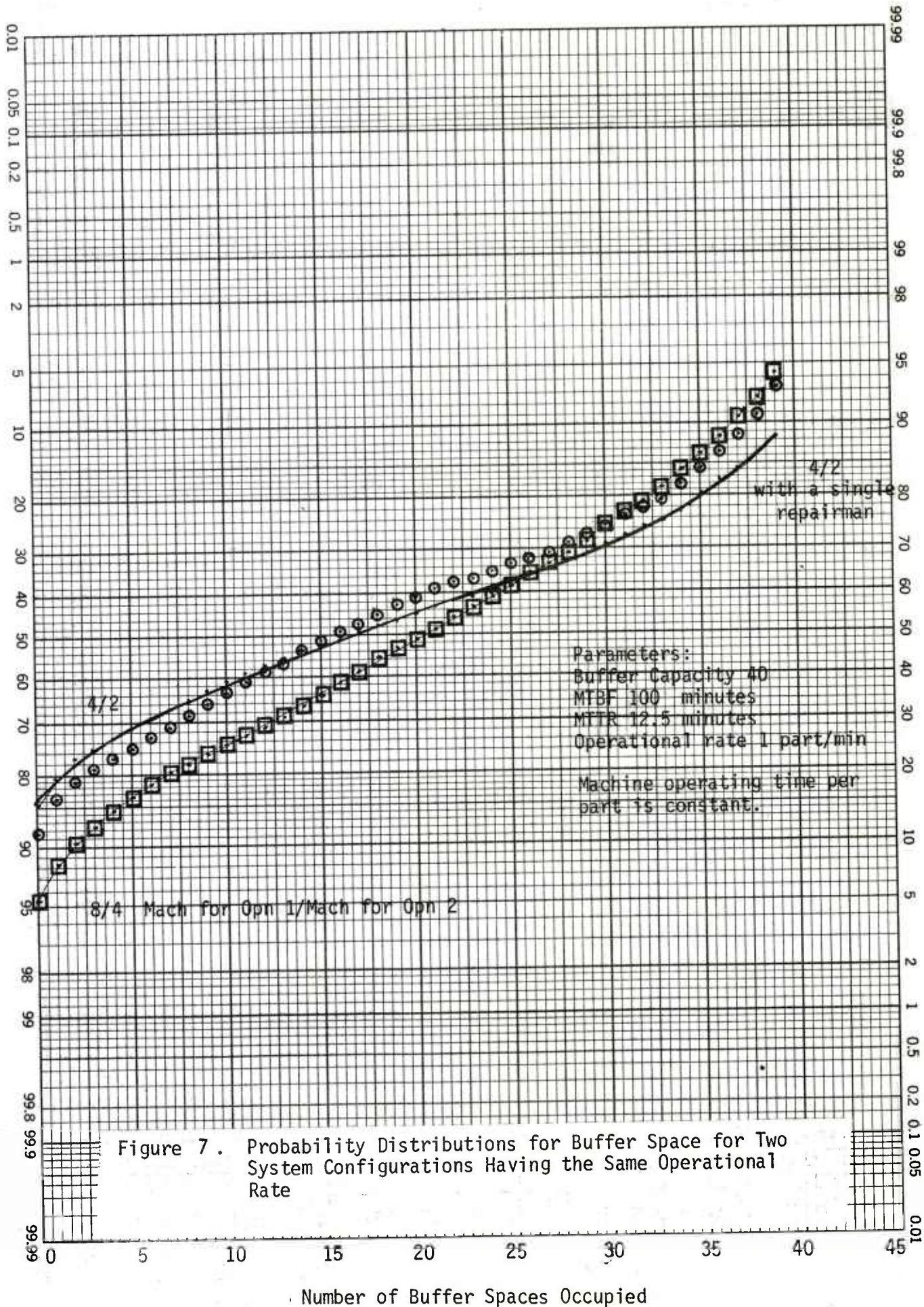
REQUIRED BUFFER CAPACITY = 32



CDF
(%)

PROBABILITY X 90 DIVISIONS
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for the slight effect of additional repairmen here is that for this system the probability that other than one repairman is needed is only about 12%. Thus, the assumption of the analysis that the system is well maintained, may not be as restrictive as it first seems.

35. Conclusions Regarding BUF.CAP

(a) The risk of exceeding the buffer capacity requirement calculated by BUF.CAP is generally quite small -- typically less than 10%. Further, the marginal productivity change, evaluated at the required capacity, is nearly a constant 0.04% per % change in buffer capacity. (b) For a large -- say, >80 spaces -- buffer, BUF.CAP executes faster than GS.BUF. This may be a consideration for execution on small computers. Actually, the execution time of BUF.CAP depends on the number of machines at each operation, not on the required buffer capacity as such. (c) Unless the number of repairmen is at least equal to the total number of machines, there is a finite probability that machines must queue for repairs. If this happens, machine availability is not equal to intrinsic availability (A) and the expected number of machines is not equal to $N \cdot A$. However, this situation is not as restrictive as it may seem. Both simulation and BUF.CAP show that the probability of exceeding a given buffer size decreases as the number of machines in each operation increases, at constant thruput. Under certain conditions the probability distribution of buffer occupancy is nearly unchanged by an increase in the number of repairmen. This occurs at a value of number of repairmen M such that $\text{prob}(\text{busy repairmen} \leq M) > 0.9$, approximately.

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ANNEXES

COMPUTER SOURCE PROGRAMS

The program listings in these annexes are written in SIMSCRIPT 2.5. They do not employ any features unique to the PRIME 550 minicomputer on which they were run. Cross-reference lists are included with the program statements to facilitate the identification of variable type and locations within the program. Since SIMSCRIPT is a language very English-like, programs can be followed easily without a flow chart. Therefore, no such diagrams are included. However, the major program blocks are announced via comments, which are distinguished from executable code by stating with double quote marks.

Potential users of these programs who do not have SIMSCRIPT compilers but do have FORTRAN compilers are assured that conversion to FORTRAN is straightforward. The code in FORTRAN is not much longer than that in SIMSCRIPT. If a FORTRAN version is implemented on a 32 bit (or less) machine, it is recommended that double-precision arithmetic be used. This is necessary to avoid truncation error when inverting large matrices.

Both main programs were designed for running interactively. Inputs are assigned following prompting messages sent to the terminal.

ANNEX 1

PROGRAM GS.BUF

A sketch of the method used to calculate the stochastic steady state of a simple production system is provided in the body of this memorandum under "Methodology Overview." Detailed system equations are derived in the section "Derivation of Equations for GS.BUF."

Input data is provided from the terminal in response to prompting messages such as "Input the operating rate for the 1st operation in parts per minute." Output is sent directly to the terminal for display. Since this output is often lengthy, it is recommended that a COMO file be established to display or print it.

Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, TRACE3
 CACI SIMSCRIPT 11.5 for PRIME Systems, Release 2.1
 27 JUL 1983 15:46:14

```

1  MAIN **FOR GS.BUF
2  DEFINE BUF.CAP, 1, J, K, L, M, N: NS AS INTEGER VARIABLES
3  DEFINE **STATE DESCRIPTION ARRAY, SDA AS AN INTEGER, 2-DIMENSIONAL ARRAY
4  **FIRST INDEX OF SDA REFERS TO THE STATE NUMBER AND THE SECOND INDEX REFERS
5  **TO THE SYSTEM ELEMENT NUMBER--1, FOR OPERATION 1; 2, FOR THE BUFFER;
6  **AND 3, FOR OPERATION 2.
7  **
8  DEFINE BV, PV, AND BUF.STATE AS REAL, 1-DIMENSIONAL ARRAYS
9  DEFINE AM, BM, AND AMINV AS REAL, 2-DIMENSIONAL ARRAYS
10 **
11 **GET BUFFER CAPACITY FROM THE TERMINAL.
12 **
13 **
14 PRINT 1 LINE THUS
15 THE INTEGER BUFFER CAPACITY.
16 READ BUF.CAP LE 2
17 IF BUF.CAP LE 2
18 PRINT 1 LINE THUS
19 BUFFER CAPACITY IS TOO SMALL.
20 INPUT ERROR. STOP
21 OTHERWISE
22 LET M=BUF.CAP+1 **BUFFER STATES
23 LET NS=4*(BUF.CAP+2) **NUMBER OF SYSTEM STATES
24 LET N=NS-1 **DIMENSION OF TRANSITION MATRIX
25 RESERVE SDA(**) AS NS BY 3
26 RESERVE AM(**) AS N BY N
27 RESERVE BM(**) AS NS BY NS
28 RESERVE BV(**) AS N
29 RESERVE BUF.STATE(*) AS M
30 **ASSIGN PARAMETER VALUES.
31 **
32 PRINT 1 LINE THUS
33 THE OPERATING RATE FOR THE 1 ST OPERATION IN PARTS PER MINUTE.
34 READ RATE1
35 LET R1=RATE1
36 PRINT 1 LINE THUS
37 THE OPERATING RATE FOR THE 2 ND OPERATION IN PARTS PER MINUTE.
38 READ RATE2
39 LET R2=RATE2
40 PRINT 1 LINE THUS
41 THE MTBF FOR THE 1 ST OPERATION IN MINUTES.
42 READ MTBF1
43 PRINT 1 LINE THUS
44 THE MTBF FOR THE 2 ND OPERATION IN MINUTES.
45 READ MTBF2
46 PRINT 1 LINE THUS
47 THE MTTR FOR THE 1 ST OPERATION IN MINUTES.
48 READ MTTR1
49 PRINT 1 LINE THUS
50 THE MTTR FOR THE 2 ND OPERATION IN MINUTES.
51 READ MTTR2
52 **
53 **CALCULATE MARKOV PARAMETERS.
54 **
55 LET MU1=1.0/MTBF1
56 LET MU2=1.0/MTBF2

```

```
50 LET LA1=1.0/MTTR1
51 LET LA2=1.0/MTTR2
52 LET S1=1.0
53 LET S2=1.0
54 FOR I=1 TO 5 DO
55   FOR J=1 TO 3 DO
56     LET SDA(I,J)=0
57   LOOP OVER J
58   LOOP OVER I
59   LET SDA(1,1)=1
60   LET SDA(2,1)=2
61   LET SDA(2,3)=1
62   LET SDA(4,3)=2
63   LET SDA(5,1)=2
64   LET SDA(6,1)=2
65   LET SDA(6,2)=0
66   LET SDA(6,3)=2
67   LET SDA(7,1)=0
68   LET SDA(7,2)=1
69   LET SDA(8,1)=0
70   LET SDA(8,2)=1
71   LET SDA(8,3)=2
72   LET SDA(9,1)=2
73   LET SDA(9,2)=1
74   LET SDA(9,3)=0
75   LET SDA(10,1)=2
76   LET SDA(10,2)=1
77   LET SDA(10,3)=2
78   FOR K=2 TO BUF.CAP-1 DO
79     LET SDA(7+4*(K-1),1)=0
80     LET SDA(7+4*(K-1),2)=K
81     LET SDA(8+4*(K-1),3)=0
82     LET SDA(8+4*(K-1),1)=0
83     LET SDA(8+4*(K-1),2)=K
84     LET SDA(8+4*(K-1),3)=2
85     LET SDA(9+4*(K-1),1)=2
86     LET SDA(9+4*(K-1),2)=K
87     LET SDA(9+4*(K-1),3)=0
88     LET SDA(10+4*(K-1),1)=2
89     LET SDA(10+4*(K-1),2)=K
90     LET SDA(10+4*(K-1),3)=2
91   LOOP OVER BUFFER STATES
92   LET SDA(NS-5,1)=0
93   LET SDA(NS-5,2)=BUF.CAP
94   LET SDA(NS-5,3)=0
95   LET SDA(NS-4,1)=0
96   LET SDA(NS-4,2)=BUF.CAP
97   LET SDA(NS-4,3)=2
98   LET SDA(NS-3,1)=2
99   LET SDA(NS-3,2)=BUF.CAP
100  LET SDA(NS-3,3)=0
101  LET SDA(NS-2,1)=2
102  LET SDA(NS-2,2)=BUF.CAP
103  LET SDA(NS-2,3)=2
104  LET SDA(NS-1,1)=1
105  LET SDA(NS-1,2)=BUF.CAP
106  LET
```

```

107 LET SDA(NS-1,3)=0
108 LET SDA(NS,1)=1
109 LET SDA(NS,2)=BUF.CAP
110 LET SDA(NS,3)=2
111 LET AVAIL1=MTBF1/(MTBF1+MTTR1)
112 LET AVAIL2=MTBF2/(MTBF2+MTTR2)
113 **
114 **ECHO PARAMETER VALUES.
115 **
116 SKIP 2 LINES
117 PRINT 12 LINES WITH BUF.CAP,RATE1,MTBF1,MTTR1,AVAIL1,RATE2,MTBF2,MTTR2,
118 AVAIL2 THUS
PARAMETER VALUES FOR THE STEADY-STATE BUFFER MODEL

```

BUFFER CAPACITY * SPACES

OPERATION NUMBER	MACH RATE (PARTS/MIN)	MTBF (MIN)	MTTR (MIN)	AVAIL- ABILITY
1	****	****	****	*****
2	****	****	****	*****

```

119 **
120 **FILL RIGHT HAND VECTOR.
121 **
122 ** FOR I=1 TO N, LET BV(I)=0.0
123 **
124 **FILL ELEMENTS OF THE STATE TRANSITION MATRIX.
125 **
126 FOR I=1 TO NS DO
127   FOR J=1 TO NS DO
128     LET BM(I,J)=0.0
129     LOOP OVER J
130   LOOP OVER I
131     LET BM(1,1)=-LA1
132     LET BM(1,2)=MU1
133     LET BM(1,4)=R2
134     LET BM(2,1)=-LA1
135     LET BM(2,2)=-MU1+R1
136     LET BM(2,6)=S2+R2
137     LET BM(3,3)=-LA1+LA2
138     LET BM(3,4)=MU2
139     LET BM(3,5)=MU1
140     LET BM(4,3)=LA2
141     LET BM(4,4)=-LA1+MU2+R2
142     LET BM(4,6)=MU1
143     LET BM(4,8)=R2
144     LET BM(5,3)=LA1
145     LET BM(5,5)=-LA2+MU1+R1
146     LET BM(5,6)=MU2
147     LET BM(6,2)=R1
148     LET BM(6,4)=LA1
149     LET BM(6,5)=LA2
150     LET BM(6,6)=-MU1+MU2+S1+R1+S2+R2
151     LET BM(6,10)=S2+R2

```

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MAIN ROUTINE

Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, IRACE3

```

152 LET BM(7,7)=-(LA1+LA2)
153 LET BM(7,8)=MU2
154 LET BM(7,9)=MU1
155 LET BM(8,7)=LA2
156 LET BM(8,8)=-(LA1+MU2+R2)
157 LET BM(8,10)=MU1
158 LET BM(8,12)=R2
159 LET BM(9,5)=R1
160 LET BM(9,7)=BM(5,3)
161 LET BM(9,9)=BM(5,5)
162 LET BM(9,10)=BM(5,6)
163 LET BM(10,6)=S1+R1
164 LET BM(10,8)=BM(6,4)
165 LET BM(10,9)=BM(6,5)
166 LET BM(10,10)=BM(6,6)
167 LET BM(10,14)=BM(6,10)
168 FOR I=1 TO NS-5 DO
169   FOR J=1-4 TO NS DO
170     LET BM(I,J)=BM(I-4,J-4)
171   LOOP
172   LOOP **OVER COLUMNS (J)
173   LET BM(NS-4,NS-5)=LA2
174   LET BM(NS-4,NS-4)=-(LA1+MU2+R2)
175   LET BM(NS-4,NS-2)=MU1
176   LET BM(NS-3,NS-7)=R1
177   LET BM(NS-3,NS-5)=LA1
178   LET BM(NS-3,NS-3)=-(LA2+MU1+R1)
179   LET BM(NS-3,NS-2)=MU2
180   LET BM(NS-2,NS-6)=S1+R1
181   LET BM(NS-2,NS-4)=LA1
182   LET BM(NS-2,NS-3)=LA2
183   LET BM(NS-2,NS-2)=-(MU1+MU2+S1+R1+S2+R2)
184   LET BM(NS-2,NS)=R2
185   LET BM(NS-1,NS-3)=R1
186   LET BM(NS-1,NS-1)=LA2
187   LET BM(NS-1,NS)=MU2
188   LET BM(NS,NS-2)=S1+R1
189   LET BM(NS,NS-1)=LA2
190   LET BM(NS,NS)=-(MU2+R2)
191 **CHECK THAT ALL COLUMN SUMS OF B(**) ARE ZERO.
192 **
193 FOR J=1 TO N+1 DO
194   LET SUM=0.0
195   FOR I=1 TO N+1, ADD BM(I,J) TO SUM
196   IF ABS(F(SUM)) GE 1.0/10.0**8
197     PRINT 1 LINE WITH J AND SUM THUS
198     PRINT ** TH COLUMN OF STATE TRANSITION MATRIX = .....
199     FOR K=1 TO N+1 DO
200       PRINT 1 LINE WITH K, BM(K,J) THUS
201     **
202     LOOP **OVER K
203     STOP
204   OTHERWISE
205   LOOP **OVER COLUMNS OF BM(**)
206 **FOR ROWS BM(*,1) NE 0.0 PERFORM ROW SUM OPERATIONS.

```

```

207  ** FOR I=2 TO N+1 DO
208  ** IF BM(I,1) NE 0.0
209  ** LET CON=BM(I,1)
210  ** FOR J=1 TO N+1, SUBTRACT CON FROM BM(I,J)
211  ** SUBTRACT CON FROM BV(I-1)
212  ** ALWAYS
213  ** LOOP ** OVER ROWS OF BM(**)
214  **
215  ** COPY THE SUPMATRIX OF BM(**) INTO AM(**).
216  **
217  ** FOR I=1 TO N DO
218  ** FOR J=1 TO N DO
219  ** LET AM(I,J)=BM(I+1,J+1)
220  ** LOOP ** OVER J
221  ** OVER I
222  ** RELEASE BM(**)
223  ** RESERVE AMINV(**) AS N BY N
224  **
225  ** CALCULATE THE MATRIX INVERSE OF AM(**).
226  **
227  ** CALL MAT-INVERSE (AM(**),N) YIELDING AMINV(**)
228  ** RESERVE PV(**) AS N
229  **
230  ** CALCULATE THE PROBABILITY STATE VECTOR PV(**).
231  **
232  ** CALL MAT-VEC.MPY (AMINV(**), BV(**), N) YIELDING PV(**)
233  **
234  ** CALCULATE THE PROBABILITY OF THE NULL STATE.
235  **
236  ** LET P-NULL=1.0
237  ** FOR I=1 TO N, SUBTRACT PV(I) FROM P-NULL
238  **
239  ** CHECK FOR VALID PROBABILITY.
240  **
241  ** IF P-NULL < 0.0 OR P-NULL > 1.0
242  ** PRINT 1 LINE WITH P-NULL THUS
243  ** IN SS-BUF. P-NULL = .....
244  ** ERROR
245  ** OTHERWISE
246  **
247  ** GET BUFFER STATE PROBABILITIES.
248  **
249  ** FOR K=1 TO M, LET BUF-STATE(K)=0.0
250  ** ADD P-NULL TO BUF-STATE(1)
251  ** FOR J=1 TO 5, ADD PV(K) TO BUF-STATE(1)
252  ** FOR K=1 TO 4 DO
253  ** LOOP ** OVER (K) SUBSTATES
254  ** OVER (J) STATES OF THE BUFFER
255  ** FOR K=1 TO 6, ADD PV(K+4*(M-1)+1) TO BUF-STATE(M)
256  ** LET P-OPN2-WAIT=PV(1)+P-NULL
257  ** LET P-OPN1-WAIT=PV(NS-1)+PV(NS-2)
258  ** LET P-SYS-WAIT=P-OPN1-WAIT+P-OPN2-WAIT
259  ** LET P-OPN2-DN=0.0
260  ** FOR K=1 TO (NS-2)/2 DO

```


MAIN ROUTINE
Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, TRACES
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```

263 ADD PV(2*K) TO P.OPN2.DN
264 LOOP %TO GET PROB THAT 2 ND OPERATION IS BEING REPAIRED
265 LET P.NO.OP=P.OPN2.DN+P.OPN2.WAIT
266 **
267 **PRINT HEADINGS.
268 **
269 SKIP 4 LINES
270 PRINT 6 LINES THUS
PROBABILITIES OF OPN1-BUFFER-OPN2 SYSTEM STATES

```

STATE NUMBER	STATE PROB	STATE DESCRIPTION OPN1 BUF OPN2
271	*****	PRINT 1 LINE WITH P.NULL, SDA(I,1), SDA(I,2), SDA(I,3) THUS
272	*****	FOR I=1 TO N DO
273	*****	PRINT 1 LINE WITH I+1, PV(I), SDA(I+1,1), SDA(I+1,2), SDA(I+1,3) THUS
274	*****	LOOP %OVER I
275	*****	PRINT 9 LINES THUS

BUFFER STEADY-STATE OCCUPANCY PROBABILITIES

BUFFER STATE	STATE PROB	STATE DESCRIPTION
276	*****	FOR I=1 TO M DO
277	*****	PRINT 1 LINE WITH I-1, BUF.STATE(I)
278	*****	THUS
279	*****	LOOP %OVER BUFFER STATES
280	*****	PRINT 5 LINES WITH P.OPN1.WAIT, P.OPN2.WAIT, P.SYS.WAIT THUS
281	*****	PROB THAT OPERATION 1 MUST WAIT WITH BUFFER FULL
282	*****	PROB THAT OPERATION 2 MUST WAIT WITH BUFFER EMPTY
283	*****	PROB THAT AN OPN IS NOT OPERATING DUE TO BUFFER
284	*****	LET E.PROD.RATE=1.0-P.NO.OP
285	*****	LET E.PROD.RATE=RATE2*PRODUCTIVITY
286	*****	PRINT 4 LINES WITH P.OPN2.DN, P.NO.OP, PRODUCTIVITY, E.PROD.RATE THUS
287	*****	PROB THAT OPERATION 2 IS DOWN FOR REPAIRS
288	*****	LET THE SYSTEM HAS NO OUTPUT RATE
289	*****	SYSTEM PRODUCTION CAPACITY (PARTS/MINUTE)
290	*****	SKIP 4 LINES
291	*****	STOP
292	*****	END MAIN FOR SS.BUF

NAME	TYPE	MODE	LINE NUMBERS OF REFERENCES
ABS.F	ROUTINE	INTEGER	197
AM	RECURSIVE VARIABLE	DOUBLE	10
AMINV	RECURSIVE VARIABLE	DOUBLE	228
AVAIL1	RECURSIVE VARIABLE	DOUBLE	10
AVAIL2	RECURSIVE VARIABLE	DOUBLE	111
BM	RECURSIVE VARIABLE	DOUBLE	112
		DOUBLE	10
		DOUBLE	135
		DOUBLE	142
		DOUBLE	149
		DOUBLE	156
		DOUBLE	163
		DOUBLE	174
		DOUBLE	181
		DOUBLE	188
		DOUBLE	211*
		DOUBLE	212
		DOUBLE	215
		DOUBLE	220
		DOUBLE	223*
		DOUBLE	225
		DOUBLE	227
		DOUBLE	233
		DOUBLE	236
		DOUBLE	241
		DOUBLE	244
		DOUBLE	251
		DOUBLE	254*
		DOUBLE	257*
		DOUBLE	263
		DOUBLE	266
		DOUBLE	272*
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```

1 ROUTINE FOR MAT.VEC.MPY (AM, BV, NELMTS) YIELDING CV
2
3 **ROUTINE TO MULTIPLY THE SQUARE MATRIX AM, OF NELMTS BY NELMTS,
4 **BY THE VECTOR BV (NELMTS BY 1), YIELDING THE VECTOR CV (NELMTS BY 1).
5
6 DEFINE I, J, K, NELMTS AS INTEGER VARIABLES
7 DEFINE BV AND CV AS REAL, 1-DIMENSIONAL ARRAYS
8 DEFINE AM AS A REAL, 2-DIMENSIONAL ARRAY
9 RESERVE BV(1) AS NELMTS
10 RESERVE CV(1) AS NELMTS
11 FOR I=1 TO NELMTS DO
12   LET CV(I)=0.0
13   FOR K=1 TO NELMTS DO
14     FOR J=1 TO NELMTS DO
15       ADD AM(I,K)*BV(K) TO CV(I)
16   LOOP **OVER K
17 LOOP **OVER I
18 RETURN
19 END **OF ROUTINE MAT.VEC.MPY

```

C R O S S - R E F E R E N C E

NAME	TYPE	L I N E N U M B E R S O F R E F E R E N C E S		
AM	ARGUMENT	1	8	9*
BV	ARGUMENT	1	7	10*
CV	ARGUMENT	1	7	11*
I	RECURSIVE	6	12*	13
J	RECURSIVE	6	14*	15*
K	ROUTINE	1	6	15*
MAT.VEC.MPY	ROUTINE	1	6	9*
NELMTS	ARGUMENT	1	6	10

Release 2.1
27 JUL 1983 15:46:14

CACI SIMSCRIPT II.5 for PRIME Systems

Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, TRACES

```

1 ROUTINE FOR MAT.INVERSE (AM, N) YIELDING BM
2
3 **ROUTINE TO OBTAIN THE INVERSE OF THE N BY N MATRIX AM VIA THE
4 **COMPACT FORM OF THE GAUSS-JORDAN METHOD. INVERSE IS RETURNED
5 **AS BM.
6
7 DEFINE I, J, K, N AS INTEGER VARIABLES
8 DEFINE AM AND BM AS REAL, 2-DIMENSIONAL ARRAYS
9 RESERVE AM(**) AS N BY N
10 RESERVE BM(**) AS N BY N
11
12 **COPY AM INTO BM. BM IS USED FOR GAUSSIAN REDUCTION.
13
14 FOR I=1 TO N DO
15   FOR J=1 TO N DO
16     LET BM(I,J)=AM(I,J)
17   LOOP **OVER J
18 LOOP **OVER I
19 FOR I=1 TO N DO
20   LET P=BM(I,I)
21   IF P=0.0
22     PRINT 2 LINES WITH I THUS
23     IN ROUTINE MAT.INVERSE. THE **TH DIAGONAL ELEMENT IS ZERO.
24     THE MATRIX CANNOT BE INVERTED.
25     STOP
26 OTHERWISE
27   LET BM(I,I)=1.0
28   FOR J=1 TO N DO
29     LET BM(I,J)=BM(I,J)/P
30   LOOP **OVER J
31   FOR J=1 TO N DO **THE SECOND J-LOOP
32     IF J=I
33       GO TO EOJ **END OF J-LOOP
34     OTHERWISE
35       LET P=BM(J,I)
36       LET BM(J,I)=0.0
37       FOR K=1 TO N DO
38         SUBTRACT P*BM(I,K) FROM BM(J,K)
39       LOOP **OVER K
40 EOJ **LOOP **OVER J
41 RETURN
42 END **OF ROUTINE MAT.INVERSE

```


CROSS - R E F E R E N C E

NAME	TYPE	NO.	NO.	MODE	LINE NUMBERS OF REFERENCES
AM	ARGUMENT	1	(2-D)	DOUBLE	1 8 9* 16
RM	ARGUMENT	3	(2-D)	DOUBLE	1 8 10* 16
EUJ	ARGUMENT				33 36*
I	LABEL				34
J	RECURSIVE VARIABLE	1		INTEGER	38 14* 16* 19* 20*
	RECURSIVE VARIABLE	2		INTEGER	30 33 34 26* 27*
K	RECURSIVE VARIABLE	3		INTEGER	15* 16* 26*
MAT-INVERSE	ROUTINE			INTEGER	34 35*
N	ARGUMENT	2		INTEGER	7 9* 10* 15
P	RECURSIVE VARIABLE	4		DOUBLE	26 27 33 36

ANNEX 2

PROGRAM BUF.CAP

BUF.CAP calculates a recommended capacity for a buffer separating two tandem production operations. The methods employed in BUF.CAP are outlined in the body of this memorandum under "A Second Approach." The system equations for each production operation are displayed in various places depending on the number of machines (N) in a given operation. For N=1, see equation (26); for N=2, see (39, 40); and for N>2, a general form is provided in equation (44). The elements of the matrix A in (44) are given in the source program in LET statements.

Input data is provided from the terminal in response to prompting messages such as "Input the number of machines in 1st operation." To assure a system balance, the machine rates assumed in BUF.CAP, i.e. calculated endogenously, may not equal the actual or desired rates. To account for this absence of input rates, the buffer requirement from BUF.CAP must be scaled in proportion to the ratio of actual thruput to assumed thruput. The program output is sent to the terminal for display. This includes an echo of all program inputs. No output files are created. If the output is to be saved, a COMO file must be established.

```

1  **BUF*CAP
2  PREAMBLE
3  NORMALLY MODE IS REAL
4  DEFINE BINOM.DENS AS A REAL FUNCTION GIVEN 3 ARGUMENTS
5  DEFINE BINOM.DIST AS A REAL FUNCTION GIVEN 3 ARGUMENTS
6  END

```

C R O S S - R E F E R E N C E

NAME	TYPE	MODE	LINE NUMBERS OF REFERENCES
BINOM.DENS	ROUTINE	DOUBLE	4
BINOM.DIST	ROUTINE	DOUBLE	5

```

1  MAIN
2  **
3  ** DRIVER FOR BUF.CAP
4  **
5  DEFINE I, IPRINT, AND REQ.CAP AS INTEGER VARIABLES
6  DEFINE NV AS AN INTEGER, 1-DIMENSIONAL ARRAY
7  DEFINE RATEV, MTBFV, MTRV, AVAILV AS REAL, 1-DIMENSIONAL ARRAYS
8  RESERVE NV(*) AS 2, MTBFV(*), MTRV(*), AVAILV(*) AS 2
9  RESERVE RATEV(*), MTBFV(*), MTRV(*), AVAILV(*) AS 2
10 **
11 ** GET INPUT PARAMETERS FROM THE TERMINAL.
12 **
13 PRINT 1 LINE THUS
14 INPUT THE NUMBER NV(1)
15 PRINT 1 LINE THUS
16 INPUT THE NUMBER OF MACHINES IN 2 ND OPERATION.
17 READ NV(2)
18 FOR I=1 TO 2 DO
19   PRINT 1 LINE WITH I THUS
20   PRINT TIME BETWEEN FAILURES (MINUTES) FOR MACHINES OF OPERATION *.
21   INPUT THE MEAN READ MTBFV(I)
22   PRINT 1 LINE WITH I THUS
23   PRINT TIME TO REPAIR MACHINES OF OPERATION *.
24   INPUT THE MEAN READ MTRV(I)
25   LET AVAILV(I) = MTBFV(I) / (MTBFV(I) + MTRV(I))
26   LOOP ** OVER OPERATIONS
27   LET MAXA = MAX.F (AVAILV(1), AVAILV(2))
28   FOR I=1 TO 2, LET RATEV(I) = 1.0 / AVAILV(I) / NV(I) * MAXA
29   PRINT 1 LINE, THUS
30   DETAILED READ IPRINT
31   CALL BUF.CAP (NV(*), RATEV(*), MTBFV(*), MTRV(*), IPRINT) YIELDING REQ.CAP
32   REQUIRED BUFFER CAPACITY = ***
33   STOP
34   ** MAIN
35 END

```

CROSS - R E F E R E N C E

NAME	TYPE	WORD	(1-D)	MODE	LINE NUMBERS OF REFERENCES									
AVAILV	RECURSIVE VARIABLE	WORD	8	DOUBLE	7	9*	22	24*	25					
BUF-CAP	ROUTINE			INTEGER	28									
I	RECURSIVE VARIABLE	WORD	1	INTEGER	25*	17*	18	19	20	21	22*			
IPRINT	RECURSIVE VARIABLE	WORD	2	INTEGER	5	27	28							
MAX-F	ROUTINE			INTEGER	24									
MAXA	RECURSIVE VARIABLE	WORD	20	DOUBLE	24	25								
MTBFV	RECURSIVE VARIABLE	WORD	6	(1-D)	7	9*	19	22*	28					
MTTRV	RECURSIVE VARIABLE	WORD	7	(1-D)	7	9*	21	22	28					
NV	RECURSIVE VARIABLE	WORD	4	(1-D)	6	8*	14	16	25	28				
RATEV	RECURSIVE VARIABLE	WORD	5	(1-D)	7	9*	25	28						
REG-CAP	RECURSIVE VARIABLE	WORD	3	INTEGER	5	28								

Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, TRACE3
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```

1 ROUTINE FOR BUF.CAP (NV, RATEV, MTBFV, MTRV, IPRINT) YIELDING REQ.CAP
2
3 **ROUTINE TO ESTIMATE THE REQUIREMENT FOR BUFFER CAPACITY BETWEEN
4 **TWO MACHINE OPERATIONS, EACH OF WHICH MAY HAVE SEVERAL IDENTICAL
5 **MACHINES OPERATING IN PARALLEL. MACHINE MAINTENANCE QUEUES ARE
6 **DISALLOWED IN THIS MODEL.
7
8 DEFINE I, I1, IPRINT, J, J1, K, L, M, M1, M2, MAXL, N, N1, N2, REQ.CAP
9 AS INTEGER VARIABLES
10 DEFINE INDEX, NV, STATE1V, AND STATE2V AS INTEGER, 1-DIMENSIONAL ARRAYS
11 DEFINE PNV, PNV, PIV, AND BUFV AS REAL, 1-DIMENSIONAL ARRAYS
12 DEFINE RATEV, MTBFV, MTRV, AVAILV AS REAL, 1-DIMENSIONAL ARRAYS
13 RESERVE NV(*), RATEV(*), MTBFV(*), MTRV(*), AVAILV(*) AS 2
14 DEFINE PDN.ARRAY AS A REAL, 2-DIMENSIONAL ARRAY
15 RESERVE PDN.ARRAY(*,*) AS 2 BY *
16 LET N1=NV(1)
17 LET N2=NV(2)
18 RESERVE PNV(*) AS N1
19 RESERVE PDN.ARRAY(1,*) AS N1+1
20 RESERVE PDN.ARRAY(2,*) AS N2+1
21
22 **ASSIGN PARAMETERS.
23
24 LET DELTAT=0.5 ** (MINUTE) FOR TIME STEP
25 LET FACTOR=4.0 ** TIMES THE LARGEST MTR
26 LET STAT.CONF=1.0-EXP.F(-FACTOR)
27 LET M=TRUNC.F(FACTOR*MAX.F(MTRV(1), MTRV(2))/DELTAT)
28 RESERVE PNV(*) AS M
29 LET M1=TRUNC.F(FACTOR*MTRV(1)/DELTAT)
30 LET M2=TRUNC.F(FACTOR*MTRV(2)/DELTAT)
31 LET MAXL=(NV(1)+1)*(NV(2)+1)
32 RESERVE PIV(*) AND BUFV(*) AS MAXL
33 RESERVE INDEX(*) AND STATE1V(*) AND STATE2V(*) AS MAXL
34 FOR K=1 TO MAXL, LET INDEX(L)=L
35 FOR K=1 TO 2, LET AVAILV(K)=MTBFV(K)/(MTBFV(K)+MTRV(K))
36 LET A1=AVAILV(1)
37 LET A2=AVAILV(2)
38 LET RATE=RATEV(1)
39 LET RATE=RATEV(2)
40 LET MTBF=MTBFV(1)
41 LET MTBF=MTBFV(2)
42
43 **ECHO INPUT DATA.
44
45 SKIP 4 LINES
46 PRINT 10 LINES WITH N1, N2, RATEV(1), RATEV(2), MTBFV(1), MTBFV(2),
47 MTRV(1), MTRV(2), 100.0*A1, 100.0*A2, M+DELTAT, STAT.CONF
48
49 MACHINE INPUT DATA FOR BUFFER CAPACITY CALCULATION
50 OPERATION 1 OPERATION 2

```

NO MACHINES		***	***	***	PARTS/MIN
MACH RATE	---	***	***	***	MINUTES
MTBF	---	***	***	***	MINUTES
MTR	---	***	***	***	PERCENT
AVAILABILITY	---	***	***	***	
REPAIR TIME (MINUTES)		***		ASSOCIATED STAT CONFIDENCE	

ROUTINE BUF,CAP CACI SIMSCRIPT II.5 for PRIME Systems, Release 2.1
Options = SEQUENCE,ID,SUBCHK,XREF,NOEXPLIST,TRACE3 27 JUN 1983 12:05:18

```

48  SKIP 4 LINES
49  FOR I1=1 TO N1+1 DO
50    LET I=I1-1
51  **
52  **INITIALIZE PROBABILITY STATE VECTOR FOR 1 ST OPERATION.
53  **
54  FOR K=1 TO N1, LET PNV(K)=0.0
55  IF I NE 0
56    LET PNV(I)=1.0
57  ALWAYS
58  CALL OPN.DYN GIVEN PNV(*),MTBF,MTTR,RATE,DELTAT,IPRINT,M
59  YIELDING PDNV(*)
60  LET PDN.ARRAY(1,I1)=PDNV(M)
61  LOOP **OVER STATES OF 1 ST OPERATION
62  RELEASE PNV(*)
63  RESERVE PNV(*) AS N2
64  LET RATE=RATEV(2)
65  LET MTBF=MTBFV(2)
66  LET MTTR=MTTRV(2)
67  FOR J1=1 TO N2+1 DO
68    LET J=J1-1
69  **
70  **INITIALIZE PROBABILITY STATE VECTOR FOR 2 ND OPERATION.
71  **
72  FOR K=1 TO N2, LET PNV(K)=0.0
73  IF J NE 0
74    LET PNV(J)=1.0
75  ALWAYS
76  CALL OPN.CYN GIVEN PNV(*),MTBF,MTTR,RATE,DELTAT,IPRINT,M
77  YIELDING PDNV(*)
78  LET PDN.ARRAY(2,J1)=PDNV(M)
79  LOOP **OVER STATES OF THE 2 ND OPERATION
80  **
81  **
82  **COMBINE OPERATIONAL STATES TO PRODUCE SYSTEM STATES.
83  **
84  LET AVGB=0.0
85  LET VARIB=0.0
86  FOR I1=1 TO N1+1 DO
87    LET I=I1-1
88    LET PROB.STATE1=BINOM.DENS(A1,N1,I)
89    FOR J1=1 TO N2+1 DO
90      LET J=J1-1
91      LET PROB.STATE2=BINOM.DENS(A2,N2,J)
92      LET L=J1+(N2+1)*I
93      LET STATE1V(L)=I
94      LET STATE2V(L)=J
95      LET PIV(L)=PROB.STATE1*PROB.STATE2
96      LET BUFV(L)=PDN.ARRAY(1,I1)-PDN.ARRAY(2,J1)
97      ADD BUFV(L)*BUFV(L) TO AVGB
98      ADD PIV(L)*BUFV(L)**2 TO VARIB
99  LOOP **OVER (J) 2 ND OPERATION STATES
100  LOOP **OVER (I) 1 ST OPERATION STATES
101  **
102  **RANK ORDER THE VALUES IN BUFV(*).
103  **
104  CALL RANK.ORDER GIVEN BUFV(*) AND INDEX(*)

```

ROUTINE BUF.CAP
Options = SEQUENCE, ID, SUBCHK, XREF, NDEXPLIST, TRACE3
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27 JUN 1983 12:05:18

```

105 LET SDB=SQRT.F(VARIG-AVGB**2)
106 LET REG.CAP=TRUNC.F(BUFV(MAXL)-BUFV(1))+4.0*SDB+0.5)
107 **
108 **PRINT HEADINGS.
109 **
110 **
111 ** SKIP 2 LINES
112 ** PRINT 7 LINES WITH M*DELTAT THUS
113 ** EXPECTED BUFFER REQUIREMENTS ORDERED OVER ALL SYSTEM STATES
114 ** BASED ON A REPAIR TIME LAG OF ***. MINUTES

```

ORDER INDEX	BUFFER SPACES	DIFFER SPACES	PROB DENS	CUML PROB	OPN 1 STATE	OPN 2 STATE
112	LET CDF=0.0					
113	FOR L=1 TO MAXL DO					
114	LET PROB=PIV(INDEX(L))					
115	ADD PROB TO CDF					
116	LET I=STATE1V(INDEX(L))					
117	LET J=STATE2V(INDEX(L))					
118	LET BUF.DIFF=BUFV(L)-BUFV(1)					
119	PRINT 1 LINE WITH L, BUFV(L), BUF.DIFF, PROB, CDF, I, J THUS					
120	LOOP **DVER (L) SYSTEM STATES					
121	PRINT 2 LINES THUS					

```

122 PRINT 3 LINES WITH AVGB AND SDB THUS
123 AVERAGE BUFFER CHANGE -----
124 STD DEV BUFFER CHANGE -----
125
126 RELEASE STATE1V(*)
127 RELEASE STATE2V(*)
128 RELEASE PUNV(*)
129 RELEASE AVAILV(*)
130 RELEASE PNV(*)
131 RELEASE BUFV(*)
132 RELEASE PIV(*)
133 RELEASE INDEX(*)
134 RELEASE PDN.ARRAY(**)
135 RETURN
136 END **MAIN FOR BUF.CAP

```

CROSS - REFERENCE

NAME	TYPE	MODE	LINE NUMBERS OF REFERENCES									
A1	RECURSIVE VARIABLE	WORD	29	36	45	88						
A2	RECURSIVE VARIABLE	WORD	31	37	45	91						
AVAILV	RECURSIVE VARIABLE	WORD	21	12	13*	35						
AVGB	RECURSIVE VARIABLE	WORD	52	84	97*	105						
BINOM.DENS	ROUTINE			88	91							
BUF.CAP	ROUTINE			1								
BUF.DIFF	RECURSIVE VARIABLE	WORD	67	118	119	96						
BUFV	RECURSIVE VARIABLE	WORD	20	11	132*	128*						
	RECURSIVE VARIABLE	WORD	63	118*	119	128*						
CDF	RECURSIVE VARIABLE	WORD	23	112	115*	119						
DELTAT	RECURSIVE VARIABLE	WORD	23	24	27	29						
EXP.F	ROUTINE			26	26	27						
FACTOR	RECURSIVE VARIABLE	WORD	25	25	50	55						
I	RECURSIVE VARIABLE	WORD	1	8	116	119						
I1	RECURSIVE VARIABLE	WORD	2	8	49*	60						
INDEX	RECURSIVE VARIABLE	WORD	14	10*	33*	104						
IPRINT	ARGUMENT	NO.	5	1	8	76						
J	RECURSIVE VARIABLE	WORD	3	8	68	74						
J1	RECURSIVE VARIABLE	WORD	4	117	119	68						
K	RECURSIVE VARIABLE	WORD	5	96	67*	78						
L	RECURSIVE VARIABLE	WORD	6	8	35*	72*						
	RECURSIVE VARIABLE	WORD	6	97*	34*	93						
	RECURSIVE VARIABLE	WORD	7	119*	98*	114						
M	RECURSIVE VARIABLE	WORD	7	8	27	45						
M1	RECURSIVE VARIABLE	WORD	8	78	11	28						
M2	RECURSIVE VARIABLE	WORD	9	8	11	29						
MAX.F	ROUTINE			27	30							
MAXL	RECURSIVE VARIABLE	WORD	10	8	31	33*						
MIRF	RECURSIVE VARIABLE	WORD	35	39	58	76						
MIRFV	RECURSIVE VARIABLE	WORD	3	1	12	35*						
MIRV	ARGUMENT	NO.	37	40	58	76						
MIRV	RECURSIVE VARIABLE	WORD	4	1	13*	27*						
N	RECURSIVE VARIABLE	WORD	11	8	16	19						
N1	RECURSIVE VARIABLE	WORD	12	8	18	19						
N2	RECURSIVE VARIABLE	WORD	13	86	88	45						
NV	ARGUMENT	NO.	1	89	91	63						
OPN.DYN	ROUTINE			58	10	16						
PDN.ARRAY	RECURSIVE VARIABLE	WORD	22	14	15*	20*						
PONV	RECURSIVE VARIABLE	WORD	18	131*	19*	78						
PIV	RECURSIVE VARIABLE	WORD	19	11	28*	60						
PNV	RECURSIVE VARIABLE	WORD	17	11	32*	97						
PROB	RECURSIVE VARIABLE	WORD	65	114	18*	56						

PROB.STATE1	RECURSIVE VARIABLE	WORD	56		DOUBLE	88	95		
PROB.STATE2	RECURSIVE VARIABLE	WORD	58		DOUBLE	91	95		
RANK.ORDER	ROUTINE				INTEGER	104			
RATEV	RECURSIVE VARIABLE	WORD	33	(1-D)	DOUBLE	38	58	64	76
REQ.CAP	ARGUMENT	NO.	2		DOUBLE	1	12	13*	38
SDR.F	RECURSIVE VARIABLE	WORD	60		INTEGER	1	18	106	45*
STAT.CONF	ROUTINE				DOUBLE	105	106	122	
STATE1V	RECURSIVE VARIABLE	WORD	27	(1-D)	DOUBLE	26	45	93	116
STATE2V	RECURSIVE VARIABLE	WORD	15	(1-D)	INTEGER	10	33*	94	117
TRUNC.F	ROUTINE				INTEGER	27	33*	30	124*
UIB.W	IMPLIED SUBSCRIPT	SYS	4		INTEGER	44	29	110	
VARIB	RECURSIVE VARIABLE	WORD	54		DOUBLE	85	48	105	

Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, TRACE3
 CACI SIMSCRIPT II.5 for PRIME Systems, Release 2.1 27 JUN 1983 12:05:18

```

1 ROUTINE FOR OPN.DYN GIVEN PNV, MTBF, MTR, RATE, DELTAT, IPRINT, M
2 YIELDING PNV
3
4 **PROGRAM SOLVES THE SET OF DIFFERENTIAL EQUATIONS WHICH CHARACTERIZE THE
5 **PRODUCTION OF A MANUFACTURING OPERATION CONSISTING OF N IDENTICAL MACHINES
6 **OPERATING IN PARALLEL. THE INITIAL STATE OF THIS SYSTEM IS REPRESENTED
7 **BY THE VALUES OF THE ELEMENTS OF THE N-DIMENSIONAL PROBABILITY VECTOR
8 **PNV(*), WHERE THE KTH ELEMENT IS THE PROBABILITY THAT K MACHINES ARE
9 **OPERATING. THE MEAN TIME BETWEEN FAILURES OF THIS TYPE IS
10 **ENTERED AS MTBF. THE MEAN TIME TO REPAIR A MACHINE IS ENTERED AS MTR.
11 **THE OPERATING PRODUCTION RATE OF EACH MACHINE IS ENTERED AS RATE.
12 **GENERALLY, THE TIME UNITS FOUND CONVENIENT ARE MINUTES. TO ACCOMMODATE
13 **A VARIETY OF SYSTEMS, THE PROGRAM USER MUST SUPPLY THE INTEGRATION TIME
14 **STEP IN COMPATIBLE TIME UNITS. IF A PRINTOUT OF THE TRAJECTORY OF THE
15 **STATE VECTOR IS DESIRED, THE INTEGER SWITCH IS SET TO 1. THE PROGRAM
16 **RETURNS M VALUES OF THE EXPECTED CUMULATIVE PRODUCTION. IN PDNV(*), FROM
17 **THIS SYSTEM WHEN IT OPERATES IN AN UNCONSTRAINED MANNER.
18
19 **INPUT:
20 **PNV -----
21 **
22 **
23 **MTBF -----
24 **
25 **MTR -----
26 **
27 **RATE -----
28 **
29 **DELTAT -----
30 **
31 **IPRINT -----
32 **
33 **M -----
34
35 **INTERNALLY:
36 **H -----
37 **
38 **OUTPUT:
39 **PNV(*) -----
40 **
41 DEFINE I, IPRINT, J, K, L, M, N AS INTEGER VARIABLES
42 DEFINE PNV AND PNV* AS REAL, 1-DIMENSIONAL ARRAYS
43 DEFINE PNDOT AS A REAL, 1-DIMENSIONAL ARRAY **LOCALLY
44 DEFINE AM AS A REAL, 2-DIMENSIONAL ARRAY **LOCALLY
45 RESERVE PNV(*) AS M **TIME STEPS
46 LET N=DIM.F(PNV(*))
47 RESERVE PNDOT(*) AS N **LOCALLY
48 RESERVE AM(*) AS N BY N **LOCALLY
49
50 **CALCULATE THE BIRTH- (LAMBDA) AND DEATH- (MU) RATE PARAMETERS.
51
52 LET LAMBDA=1.0/MTR
53 LET MU=1.0/MTBF
54
55 **CHECK IF TIME STEP IS REASONABLE.
56
57 IF DELTAT GE 0.1*MIN.F(MTR,MTBF)
    PRINT 3 LINES WITH DELTAT, MTR, MTR THUS

```

CACI SIMSCRIPT II.5 for PRIME Systems, Release 2.1
27 JUN 1983 12:05:18

```

ROUTINE OPN.DYN
Options = SEQUENCE,IO,SUBCHK,XREF,NOEXPLIST,TRACE3
TROUBLE IN OPN.DYN.  TIME STEP = **.* ** TOO LARGE.
MTBF = **.* **
MTTR = **.* **
      STOP
    OTHERWISE
      LET H=0.2*DELTAI **INTEGRATION STEP FOR DIFFERENTIAL EQUATIONS
      IF N > 1
        GO TO L1
      OTHERWISE **GET ANALYTIC SOLUTION FOR N=1
        LET FFREQ=LAMBDA*MU
        LET P10=PNV(1)
        LET P00=1.0-P10
        LET P0INF=MU/FFREQ
        IF IPRINT=1
          SKIP 2 LINES
          PRINT 6 LINES THUS
          TRAJECTORY OF PRODUCTION OPERATION STATE PROBABILITIES

```

TIME (MIN)	CUM PON	MACHINES OPERATING 0	Avg RATE	S.D. RATE
71	1	1		
72	1	1		
73	1	1		
74	1	1		
75	1	1		
76	1	1		
77	1	1		
78	1	1		
79	1	1		
80	1	1		
81	1	1		
82	1	1		
83	1	1		
84	1	1		
85	1	1		
86	1	1		
87	1	1		
88	1	1		
89	1	1		
90	1	1		
91	1	1		
92	1	1		
93	1	1		
94	1	1		
95	1	1		
96	1	1		
97	1	1		
98	1	1		
99	1	1		
100	1	1		
101	1	1		
102	1	1		
103	1	1		
104	1	1		

TRAJECTORY OF PRODUCTION OPERATION STATE PROBABILITIES

TIME (MIN)	CUM PDN	MACHINES OPERATING 0 1 2	AVG RATE	S.D. RATE
105	ALWAYS			
106	LET CUM=0.0			
107	LET AVGP=P10+2.0*P20			
108	FOR I=1 TO M DO			
109	LET TIME=1*DLTAT			
110	LET E1=EXP.F(-FFREQ*TIME)			
111	LET E2=EXP.F(-2.0*FFREQ*TIME)			
112	LET E3=EXP.F(-2.0*MU*TIME)			
113	LET P1=PIINF+A*E1+B*E2			
114	LET P2=P20+E3+P2INF*(1.0-E3)+A*LAMBDA/(LAMBDA-MU)*(E3-E1)+0.5*B*(E3-E2)			
115	LET P3=1.0-P1-P2			
116	LET AVGP=P1+2.0*P2			
117	LET ER2=P1+4.0*P2			
118	LET SD.RATE=RATE*SQRT.F(ER2-AVG**2)			
119	LET AVG.RATE=RATE*AVG			
120	ADD 0.5*DLTAT*(AVG+AVGP) TO CUM			
121	LET AVGP=AVG			
122	LET PDNV(1)=RATE*CUM			
123	IF IPRINT=1			
124	IF MOD.F(I,5)=0			
125	PRINT I LINE WITH TIME, PDNV(1), P0, P1, P2, AVG.RATE, SD.RATE			
126	THUS			
127	**** ALWAYS			
128	**** ALWAYS			
129	LOOP **OVER I			
130	GO TO L3			
131	** CALCULATE THE ELEMENTS OF THE STATE TRANSITION MATRIX (AM(**,**)).			
132	** L2*FOR I=1 TO N DO			
133	FOR J=1 TO N DO			
134	LET AM(I,J)=0.0			
135	LOOP **OVER J			
136	LOOP **OVER I			
137	LET CUM=N*LAMBDA			
138	LET AM(1,1)=-(CON+(N-1)*LAMBDA+MU)			
139	LET AM(1,2)=2.0*MU-CON			
140	IF N > 2			
141	FOR K=3 TO N DO			
142	LET AM(1,K)=-CON			
143	LOOP **OVER K			
144	ALWAYS			
145	IF N > 2			
146	FOR K=2 TO N-1 DO			
147	LET AM(K,K-1)=(N-K+1)*LAMBDA			
148	LET AM(K,K)=-(N-K)*LAMBDA+K*MU			
149	LET AM(K,K+1)=(K+1)*MU			
150	LOOP **OVER K			
151	ALWAYS			
152	CONTINUE TO STATE N			
153	LET AM(N,N-1)=LAMBDA			
154				

```

155 LET AM(N,N)=-N*MU
156
157 **END OF STATE TRANSITION MATRIX. GET INITIAL RATE VECTOR.
158
159 CALL MAT.VEC.MPY(AM(*,*),PNV(*),N) YIELDING PNDOT(*)
160 ADD CON TO PNDOT(1)
161
162 **PRINT HEADINGS FOR OUTPUT OF STATE VECTOR.
163
164 IF IPRINT=1
165 LET P1=0.0
166 LET P2=0.0
167 LET P3=0.0
168 LET P4=0.0
169 LET P5=0.0
170 LET P6=0.0
171 LET P7=0.0
172 LET P8=0.0
173 LET P9=0.0 **FOR PRINTING IN ROWS
174 SKIP 2 LINES
175 PRINT 7 LINES WITH N, PNV(N) THUS
TRAJECTORY OF PRODUCTION OPERATION STATE PROBABILITIES

```

INITIAL MAX STATE PROB(**) = *.****

TIME (MIN)	CUM PDN	1	2	3	4	5	6	7	8	9
176	ALWAYS									
177	**									
178	**CREATE THE EXPECTED CUMULATIVE PRODUCTION VECTOR.									
179	**									
180	LET CUM=0.0 **CUM NORMALIZED PDN FROM PREVIOUS STEP									
181	LET AVGP=0.0									
182	FOR K=1 TO N DO									
183	ADD K*PNV(K) TO AVGP									
184	LOOP **TO GET INITIAL AVERAGE NORMALIZED PRODUCTION RATE									
185	I=1 TO M **TIME STEPS** DO									
186	LET TM=I*DELTA T									
187	FOR L=1 TO 5 DO									
188	CALL MAT.VEC.CMPY (AM(**), PNV(**), N) YIELDING PNDOT(**)									
189	ADD CUM TO PNDOT(1)									
190	FOR K=1 TO N DO									
191	ADD H*PNDOT(K) TO PNV(K)									
192	LOOP **OVER (K) COMPONENTS									
193	ADD H*PNDOT(K) TO PNV(K)									
194	FOR K=1 TO N DO									
195	ADD K*PNV(K) TO AVGP									
196	LOOP **OVER COMPONENTS									
197	ADD 0.5*DELTA T*(AVGP+AVGP) TO CUM									
198	LET PNDV(I)=RATE*CUM									
199	LET AVGP=AVG									
200	IF IPRINT=1									
201	IF MOD.F(1,5)=0									
202	LET P1=PNV(1)									
203	LET P2=PNV(2)									
204										

C R O S S - R E F E R E N C E

NAME	TYPE	MODE	LINE NUMBERS OF REFERENCES			
A	RECURSIVE VARIABLE	DOUBLE	100	113	114	
A0	RECURSIVE VARIABLE	DOUBLE	96	100	101	
A1	RECURSIVE VARIABLE	DOUBLE	43	100	101	144
AM	RECURSIVE VARIABLE	DOUBLE	150	151	154	141 148 158 194 196*
AVG	RECURSIVE VARIABLE	DOUBLE	116	118	119	140 141 148 155 158 194 196*
AVG.RATE	RECURSIVE VARIABLE	DOUBLE	198	200		120
AVGP	RECURSIVE VARIABLE	DOUBLE	81	83	119	125
B	RECURSIVE VARIABLE	DOUBLE	107	120	121	183* 198 200
C	RECURSIVE VARIABLE	DOUBLE	101	113	114	
CON	RECURSIVE VARIABLE	DOUBLE	75	76		189
CUM	RECURSIVE VARIABLE	DOUBLE	139	140	141	160 189 198* 199
DELTAT	RECURSIVE VARIABLE	DOUBLE	106	120*	122	120 198* 199 73
DIM.F	ROUTINE	INTEGER	186	1	57	
E1	RECURSIVE VARIABLE	DOUBLE	45	75	79*	110 113 114
E2	RECURSIVE VARIABLE	DOUBLE	74	113	114	
E3	RECURSIVE VARIABLE	DOUBLE	111	113		
EXP.F	RECURSIVE VARIABLE	DOUBLE	112	114*		
FFREQ	RECURSIVE VARIABLE	DOUBLE	117	118		
H	ROUTINE	DOUBLE	64	110	111	112
I	RECURSIVE VARIABLE	DOUBLE	99	100*	101*	92 96 98
IPRINT	RECURSIVE VARIABLE	DOUBLE	60	191	110	
J	RECURSIVE VARIABLE	INTEGER	40	72*	76	83 108* 108*
K	RECURSIVE VARIABLE	INTEGER	109	122	124	136 136 185*
L	RECURSIVE VARIABLE	INTEGER	186	199	202	164
L0	RECURSIVE VARIABLE	INTEGER	201	241	277	102 123 230
L1	RECURSIVE VARIABLE	INTEGER	40	135*	136	150*
L2	RECURSIVE VARIABLE	INTEGER	40	143*	144	151*
L3	RECURSIVE VARIABLE	INTEGER	182*	183*	190*	196*
LAMBDA	RECURSIVE VARIABLE	INTEGER	206	187*	214	226
M	RECURSIVE VARIABLE	INTEGER	233	210	218	222
MAT.VEC.MPY	RECURSIVE VARIABLE	DOUBLE	62	89		99
MIN.F	RECURSIVE VARIABLE	DOUBLE	90	134	241	98 150
MOD.F	RECURSIVE VARIABLE	DOUBLE	88	130	275*	149
MTBF	RECURSIVE VARIABLE	DOUBLE	51	64	114*	185
MTTR	RECURSIVE VARIABLE	DOUBLE	100	101		
MU	RECURSIVE VARIABLE	DOUBLE	154	40	72	98 150 151
N	RECURSIVE VARIABLE	INTEGER	1	188	44	61 89 134 148 182
	RECURSIVE VARIABLE	INTEGER	159	1		147* 175*
	RECURSIVE VARIABLE	INTEGER	56	124	202	96 141 150 151
	RECURSIVE VARIABLE	INTEGER	78	52	56	100
	RECURSIVE VARIABLE	INTEGER	1	51	56	151
	RECURSIVE VARIABLE	INTEGER	1	64	67	98 150 151
	RECURSIVE VARIABLE	INTEGER	52	112	114	140
	RECURSIVE VARIABLE	INTEGER	101	45	46	47*
	RECURSIVE VARIABLE	INTEGER	155	139	140	61 89 134 148 182
	RECURSIVE VARIABLE	INTEGER	135	150	154*	147* 175*
	RECURSIVE VARIABLE	INTEGER	149			

Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, TRACES

```

1 ROUTINE FOR MAT.VEC.MPY (AM, BV, NELMTS) YIELDING CV
2
3 **ROUTINE TO MULTIPLY THE SQUARE MATRIX AM OF NELMTS BY NELMTS,
4 **BY THE VECTOR BV (NELMTS BY 1), YIELDING THE VECTOR CV (NELMTS BY 1).
5
6
7 DEFINE I, J, K, NELMTS AS INTEGER VARIABLES
8 DEFINE BV AND CV AS REAL, 1-DIMENSIONAL ARRAYS
9 DEFINE AM AS A REAL, 2-DIMENSIONAL ARRAY
10 RESERVE BV(*) AS NELMTS
11 RESERVE CV(*) AS NELMTS
12 FOR I=1 TO NELMTS DO
13   LET CV(I)=0.0
14   FOR K=1 TO NELMTS DO
15     FOR J=1 TO NELMTS DO
16       LET CV(I)=CV(I)+AM(I,K)*BV(K) TO CV(I)
17     LOOP **OVER K
18   LOOP **OVER I
19 RETURN
20 END **OF ROUTINE MAT.VEC.MPY

```

C R O S S - R E F E R E N C E

NAME	TYPE	NO.	1	2	3	4	5	6	7	8	9*	10	11	12	13	14	15	15*
AM	ARGUMENT	NO.	1	(2-D)	DOUBLE					1	9*	15						
ABV	ARGUMENT	NO.	2	(1-D)	DOUBLE					1	10*	15						
CCV	ARGUMENT	NO.	4	(1-D)	DOUBLE					1	11*	15						
I	RECURSIVE	WORD	1		INTEGER					6	13							
J	RECURSIVE	VARIABLE	2		INTEGER					6	13							
J	RECURSIVE	VARIABLE	3		INTEGER					6	13							
K	ROUTINE	WORD			INTEGER					14*	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
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K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
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K	ROUTINE	WORD			INTEGER					1	15*							
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K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
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K	ROUTINE	WORD			INTEGER					1	15*							
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K	ROUTINE	WORD			INTEGER					1	15*							
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K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							
K	ROUTINE	WORD			INTEGER					1	15*							

Options = SEQUENCE, ID, SUBCHK, XREF, NOEXPLIST, TRACE3

```

1  FUNCTION FOR BINOM.DENS (P, N, K)
2
3  **FUNCTION CALCULATES THE PROBABILITY DENSITY OF THE BINOMIAL DISTRIBUTION
4  **WITH (AVG) PROBABILITY PARAMETER P, WITH SAMPLE SIZE PARAMETER N, AND
5  **WITH INTEGER ARGUMENT K.
6
7  DEFINE I, K, N AS INTEGER VARIABLES
8
9  LET Q=1.0-P
10 LET BPF=Q**N
11 IF K LE 0
12   RETURN WITH BPF
13 OTHERWISE
14   FOR I=1 TO K DO
15     LET BPF=P*(N-I+1)/(I*Q)*BPF
16   LOOP **OVER I
17 END **BINOM.DENS

```

C R O S S - R E F E R E N C E

NAME	TYPE	MODE	LINE NUMBERS OF REFERENCES	
BINOM.DENS	ROUTINE	DOUBLE	1	14*
BPF	RECURSIVE VARIABLE	DOUBLE	9	14*
I	RECURSIVE VARIABLE	INTEGER	7	13*
K	RECURSIVE VARIABLE	INTEGER	1	10
N	RECURSIVE VARIABLE	INTEGER	1	9
P	RECURSIVE VARIABLE	DOUBLE	1	14
Q	RECURSIVE VARIABLE	DOUBLE	8	14

Options = SEQUENCE, I, SUBCHK, XREF, NOEXPLIST, TRACES
 CACI SIMSCRIPT II.5 for PRIME Systems, Release 2.1
 27 JUN 1983 12:05:18

```

1 FUNCTION FOR BINOM.DIST (P, N, K)
2 **
3 **FUNCTION PRODUCES THE CUMULATIVE BINOMIAL DISTRIBUTION WITH PROBABILITY
4 **PARAMETER P, SAMPLE-SIZE PARAMETER N, AND INTEGER ARGUMENT K.
5 **
6 DEFINE I, K, N AS INTEGER VARIABLES
7 LET Q=1.0-P
8 LET FX=Q**N
9 LET CDF=FX
10 IF K LE 0
11 RETURN WITH FX
12 OTHERWISE
13 FOR J=1 TO K DO
14 LET FX=P*(N-I+1)/(I*Q)*FX
15 ADD FX TO CDF
16 LOOP **OVER I
17 RETURN WITH CDF
18 END **BINOM.DIST
    
```

C R O S S - R E F E R E N C E

NAME	TYPE	MODE	LINE NUMBERS OF REFERENCES	
BINOM.DIST	ROUTINE	DOUBLE	1	15*
CDF	RECURSIVE VARIABLE	DOUBLE	9	11
FX	RECURSIVE VARIABLE	DOUBLE	8	14*
I	RECURSIVE VARIABLE	DOUBLE	6	13*
K	RECURSIVE VARIABLE	DOUBLE	1	10
N	RECURSIVE VARIABLE	DOUBLE	1	8
P	RECURSIVE VARIABLE	DOUBLE	1	14
Q	RECURSIVE VARIABLE	DOUBLE	7	14

```

1 ROUTINE TO RANK ORDER GIVEN XV AND INDEX
2 **
3 **ROUTINE ACCEPTS THE VALUES IN THE VECTOR XV, OF DIMENSION N,
4 **AND RETURNS THE VALUES ORDERED IN ASCENDING ORDER.
5 **
6 DEFINE I, II, J, AND N AS INTEGER VARIABLES
7 DEFINE INDEX AS AN INTEGER, 1-DIMENSIONAL ARRAY
8 DEFINE XV AS A 1-DIMENSIONAL ARRAY
9 LET N=DIM F(XV(*))
10 FOR I=1 TO N-1 DO
11   FOR II=I+1 TO N DO
12     IF XV(I) LE XV(II)
13       GO TO S
14     OTHERWISE **SWAP VALUES
15     LET HOLD=XV(II)
16     LET XV(II)=XV(I)
17     LET XV(I)=HOLD
18     LET J=INDEX(I)
19     LET INDEX(I)=INDEX(II)
20     LET INDEX(II)=J
21 **S*LOOP **OVER I
22 LOOP **OVER I
23 END **OF ROUTINE RANK ORDER

```

CROSS - REFERENCE

NAME	TYPE	WORD	MODE	LINE NUMBERS OF REFERENCES				
DIM, F	ROUTINE	WORD	INTEGER	9				
HOLD	RECURSIVE VARIABLE	WORD	DOUBLE	15	17			18
I	RECURSIVE VARIABLE	WORD	INTEGER	19	10*	11	12	16
II	RECURSIVE VARIABLE	WORD	INTEGER	6	11*	12	15	17
J	ARGUMENT	NO.	(1-D)	1	17	18	19	20
N	RECURSIVE VARIABLE	WORD	INTEGER	6	18	20	19*	20
RANK ORDER	ROUTINE	WORD	INTEGER	6	9	10	11	
S	ROUTINE	NO.	INTEGER	13	21			
XV	ARGUMENT	NO.	(1-D)	1	8	9	12*	15
								16*
								17